

Fast differentiable sorting and ranking



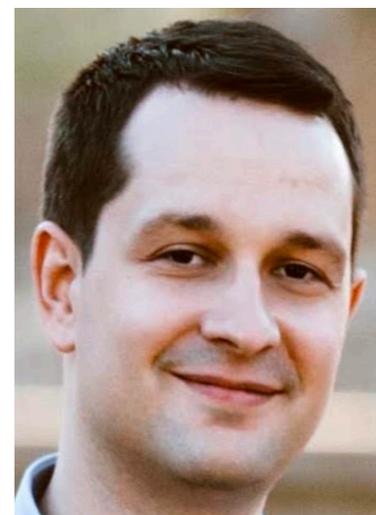
M. Blondel



O. Teboul

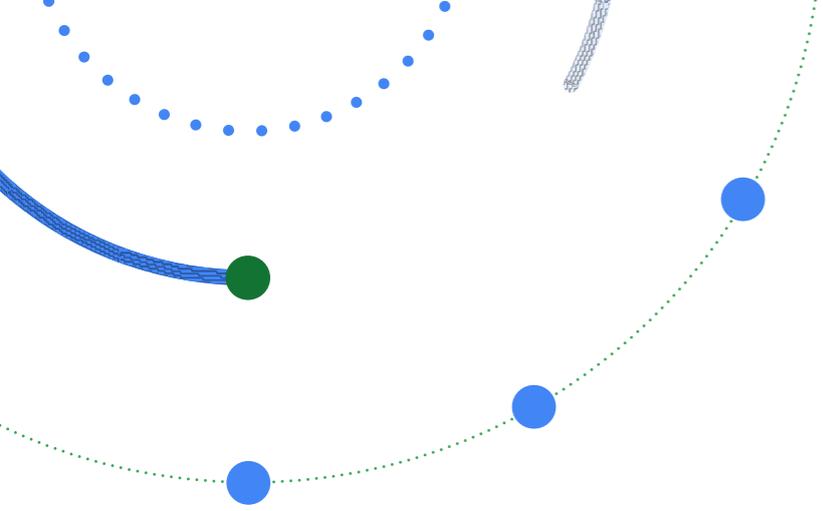


Q. Berthet



J. Djolonga

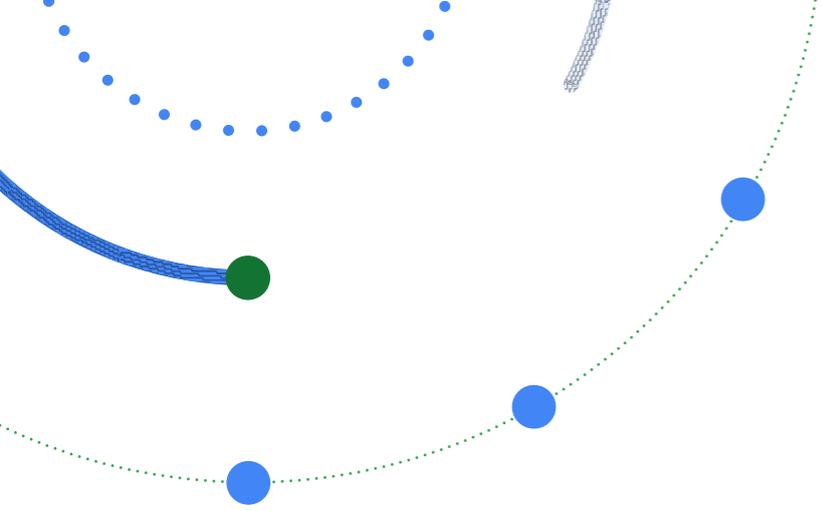
March 12th, 2020



Background

Proposed method

Experimental results



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DL as Differentiable Programming

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Deep learning increasingly synonymous with differentiable programming



Yann LeCun, 2018

“People are now building a **new kind of software** by assembling networks of parameterized **functional blocks** (including loops and conditionals) and by **training** them from examples using some form of gradient-based optimization.”

DL as Differentiable Programming

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Many computer programming operations remain **poorly differentiable**

In this work, we focus on **sorting** and **ranking**.

Sorting as subroutine in ML

k-NN

- (1) select neighbours
- (2) majority vote

Trimmed

regression

ignore large errors

Classifiers

select top-*k* activations

MoM

estimators

Ranking / Sorting

$O(n \log n)$

Learning to rank

NDCG loss and others

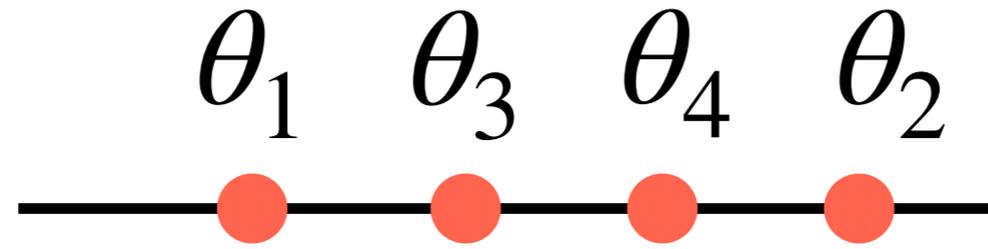
Descriptive statistics

Empirical distribution function
quantile normalization

Rank-based statistics

data viewed as ranks

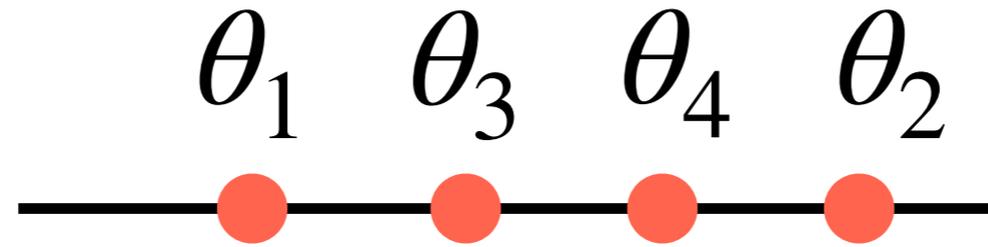
Sorting



Argsort (decending)

$$\sigma(\theta) = (2, 4, 3, 1)$$

Sorting



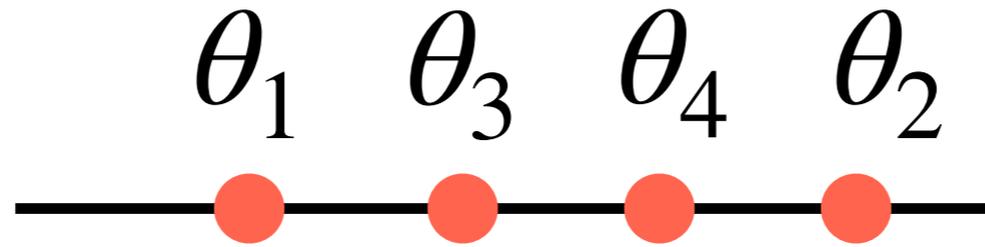
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$$s(\theta) \triangleq \theta_{\sigma(\theta)}$$

Sorting



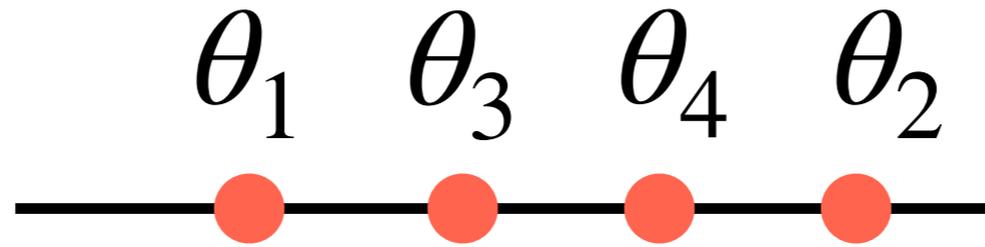
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Sorting

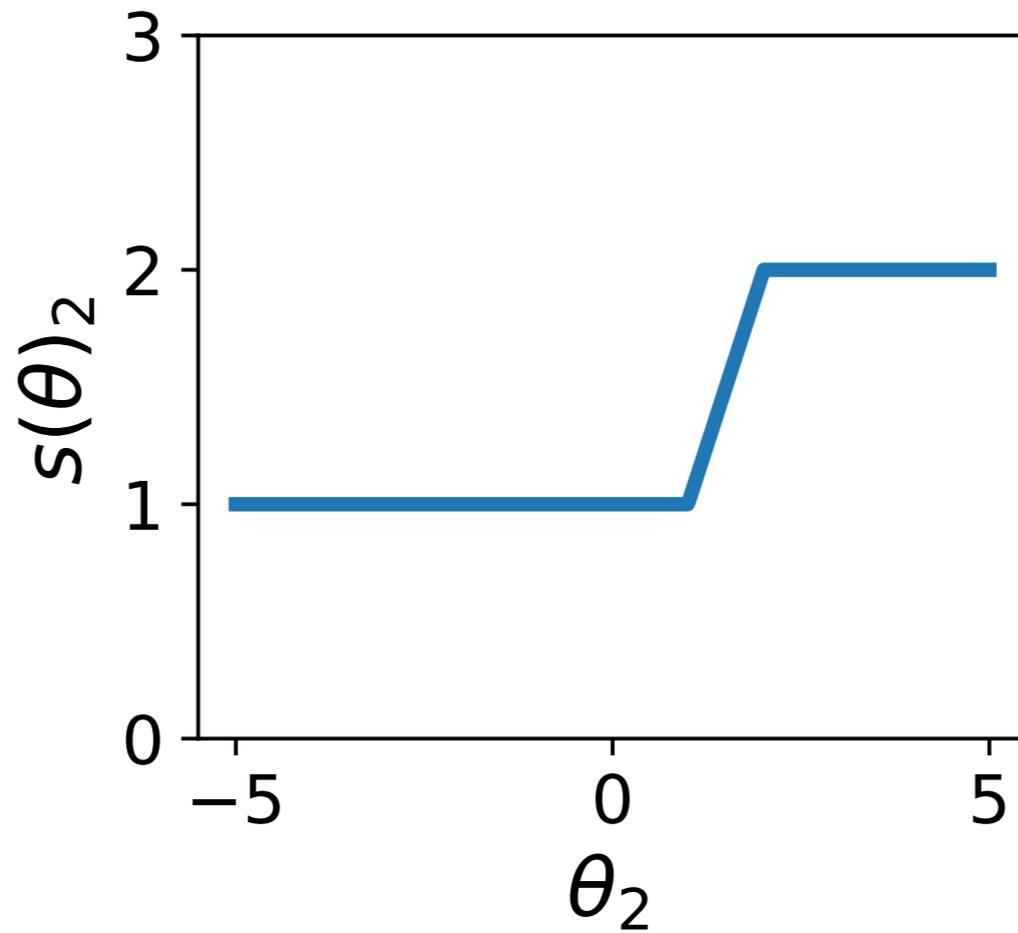


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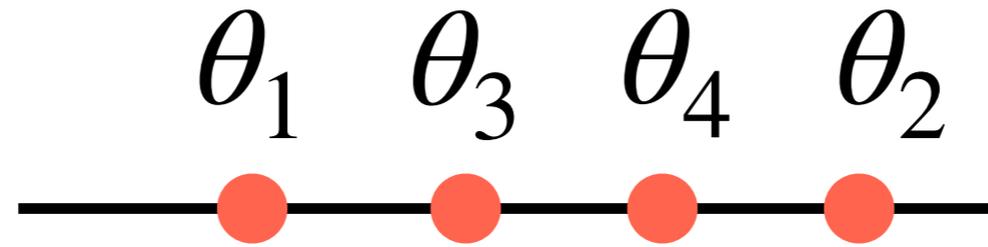


piecewise linear

induces

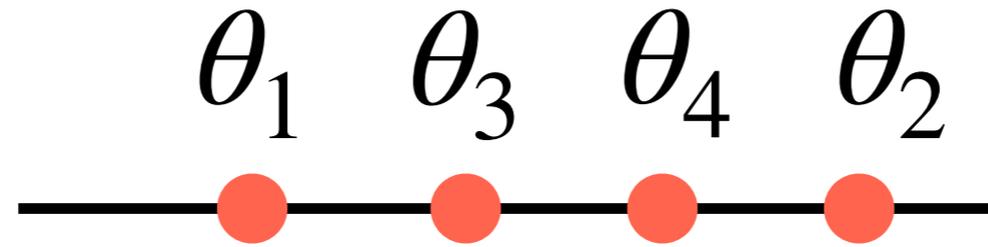
non-convexity

Ranking



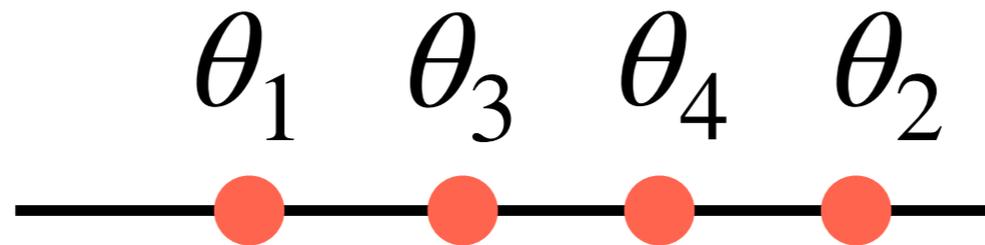
Ranks $r(\theta) \triangleq \sigma^{-1}(\theta)$

Ranking

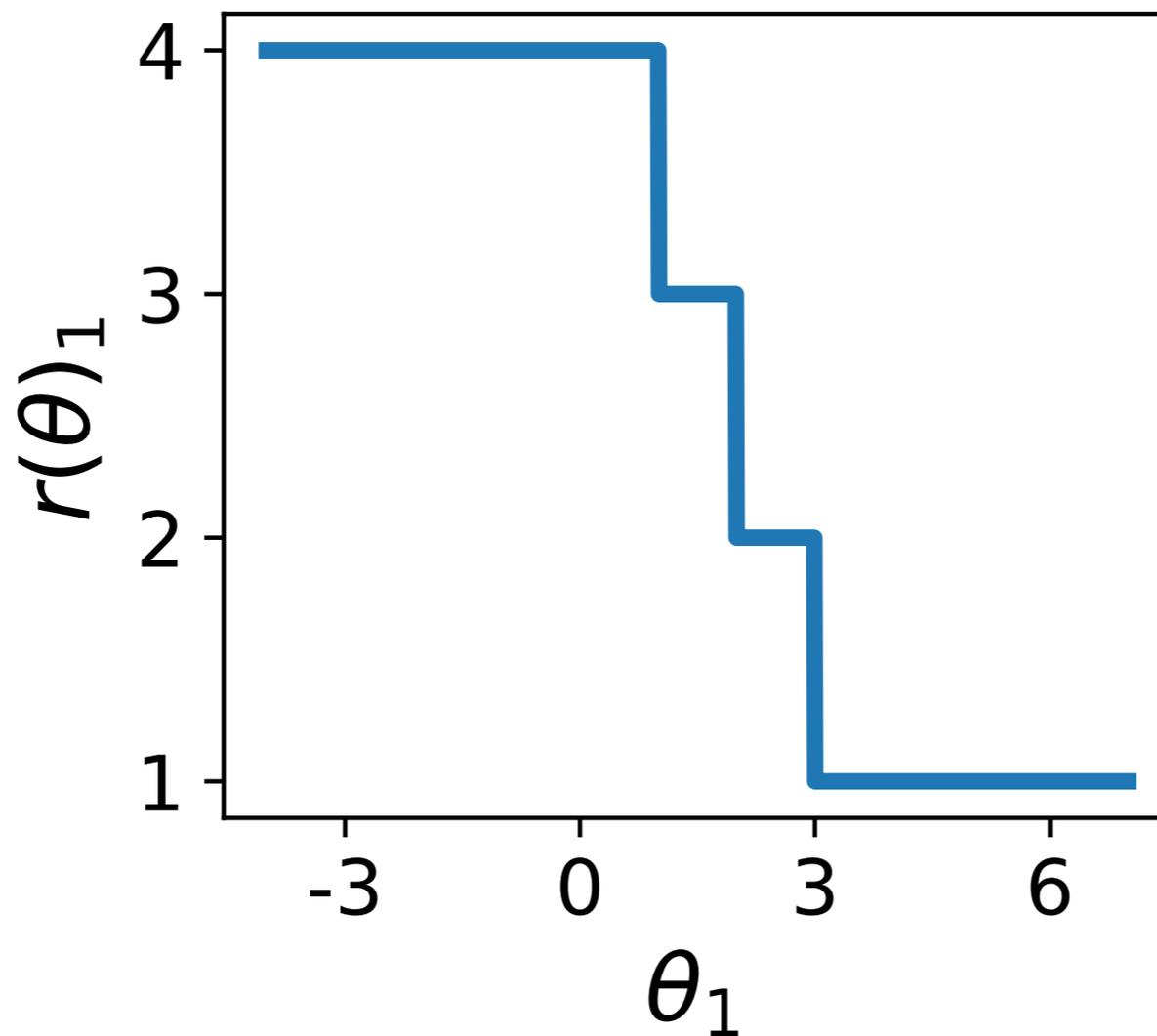


Ranks $r(\theta) \triangleq \sigma^{-1}(\theta) = (4, 1, 3, 2)$

Ranking



Ranks $r(\theta) \triangleq \sigma^{-1}(\theta) = (4, 1, 3, 2)$



discontinuous

piecewise constant

Related work on soft ranks

Soft ranks : differentiable proxies to “hard” ranks

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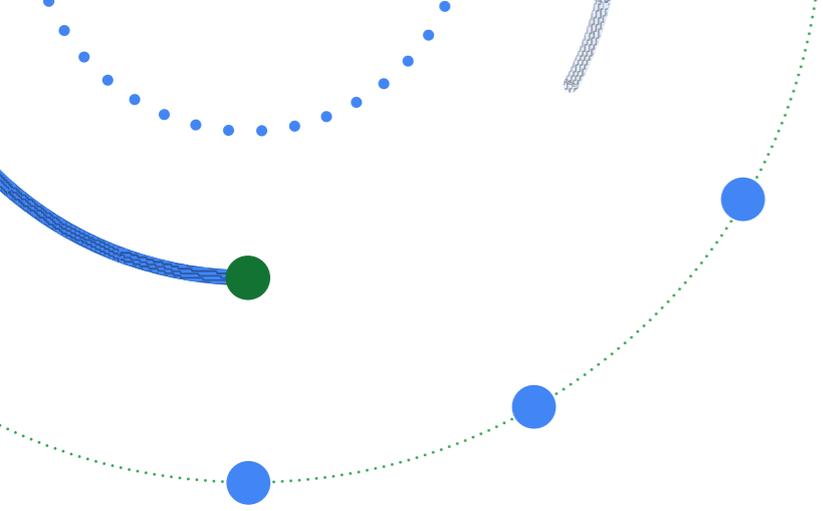
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None of these works achieves $O(n \log n)$ complexity



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Our proposal

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- Exact computation in $O(n \log n)$ time (forward pass)

Our proposal

- Differentiable (soft) relaxations of $s(\theta)$ and $r(\theta)$
- Two formulations: **L2** and Entropy regularised
- “Convexification” effect
- Exact computation in $O(n \log n)$ time (forward pass)
- Exact multiplication with the Jacobian in $O(n)$ time without unrolling (backward pass)

Strategy outline

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→ Could be challenging (argmin differentiation problem)

Strategy outline

Cuturi et al. [2019]

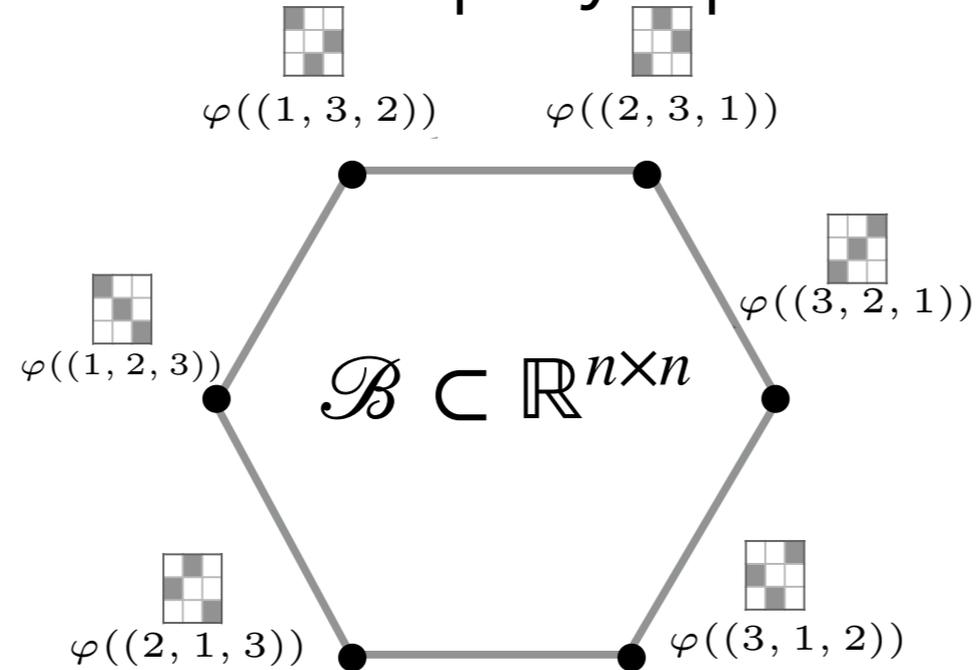
This work

Strategy outline

1. LP

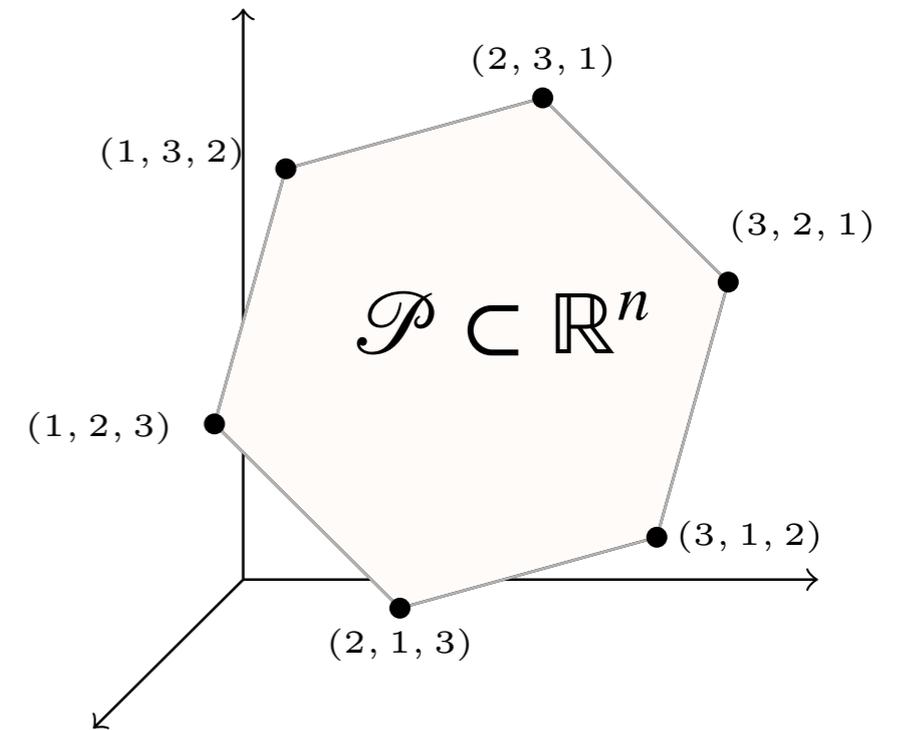
Cuturi et al. [2019]

Birkhoff polytope



This work

Permutahedron

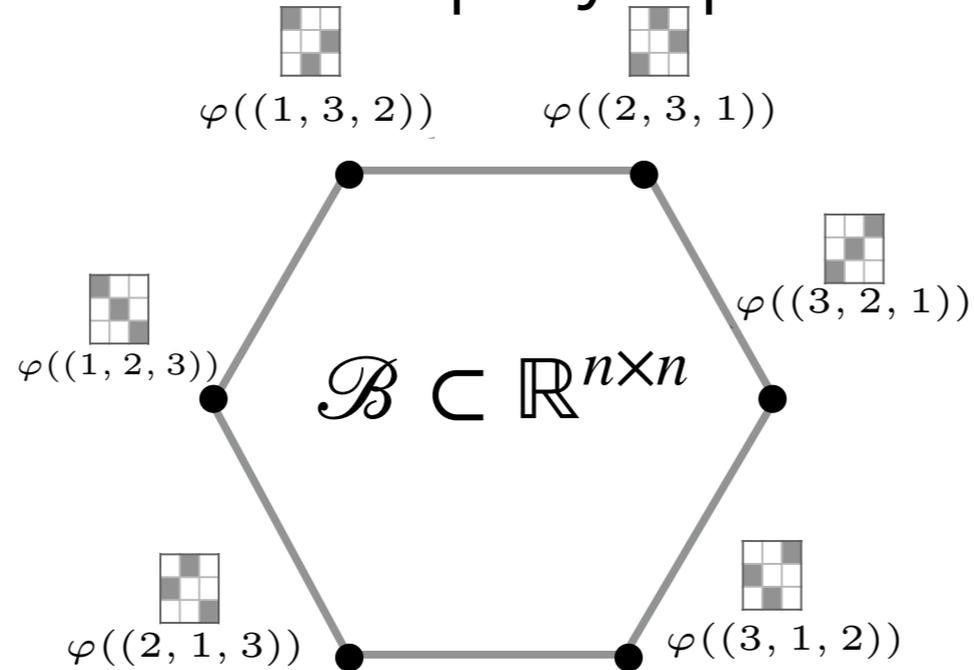


Strategy outline

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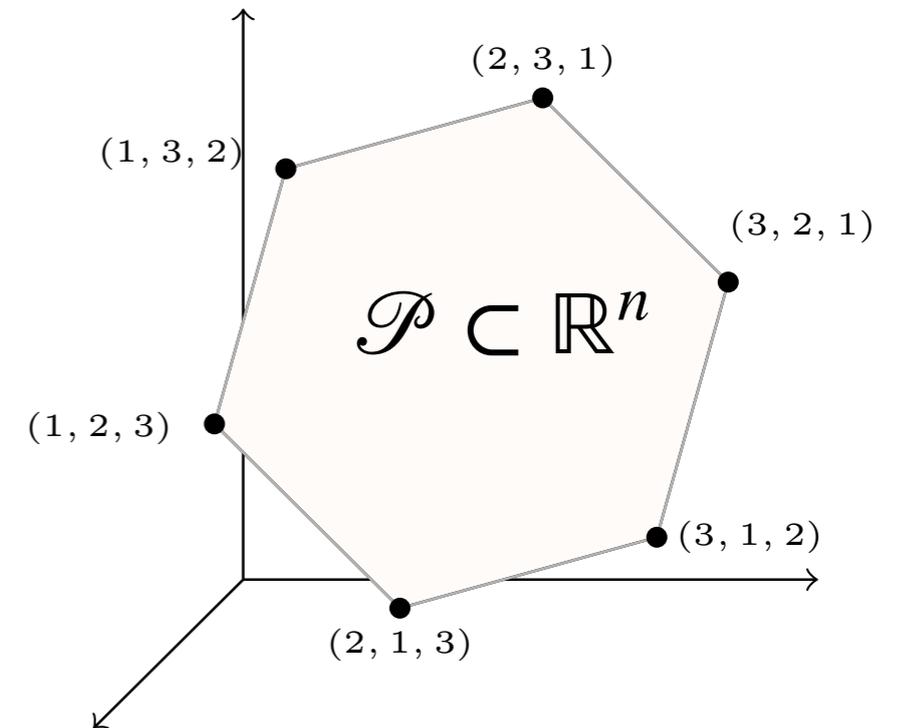
Birkhoff polytope



Entropy

This work

Permutahedron



L2 or Entropy

2. Regularization

Strategy outline

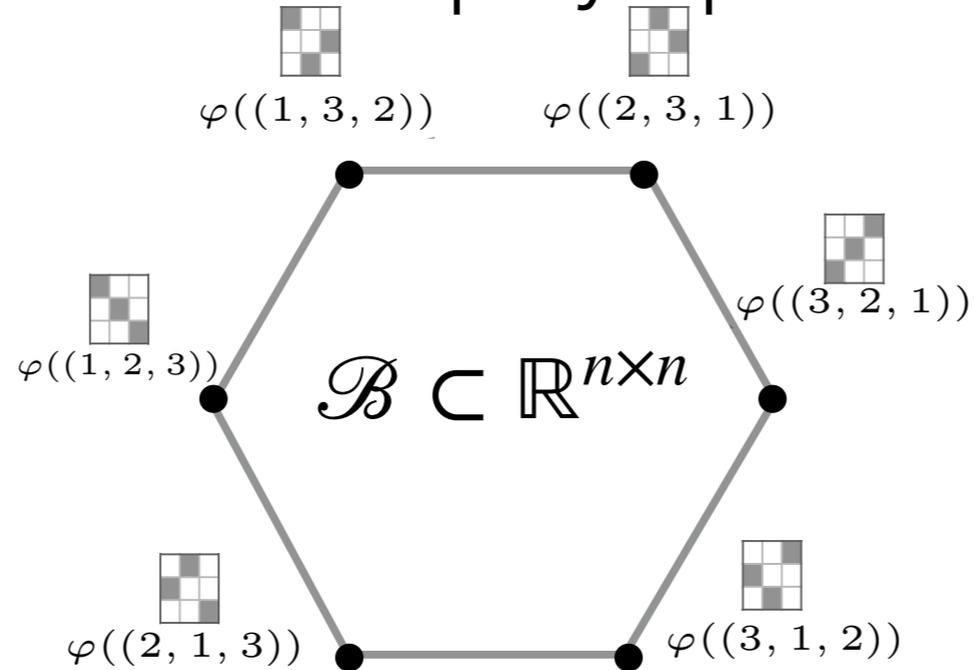
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2. Regularization

3. Computation

Cuturi et al. [2019]

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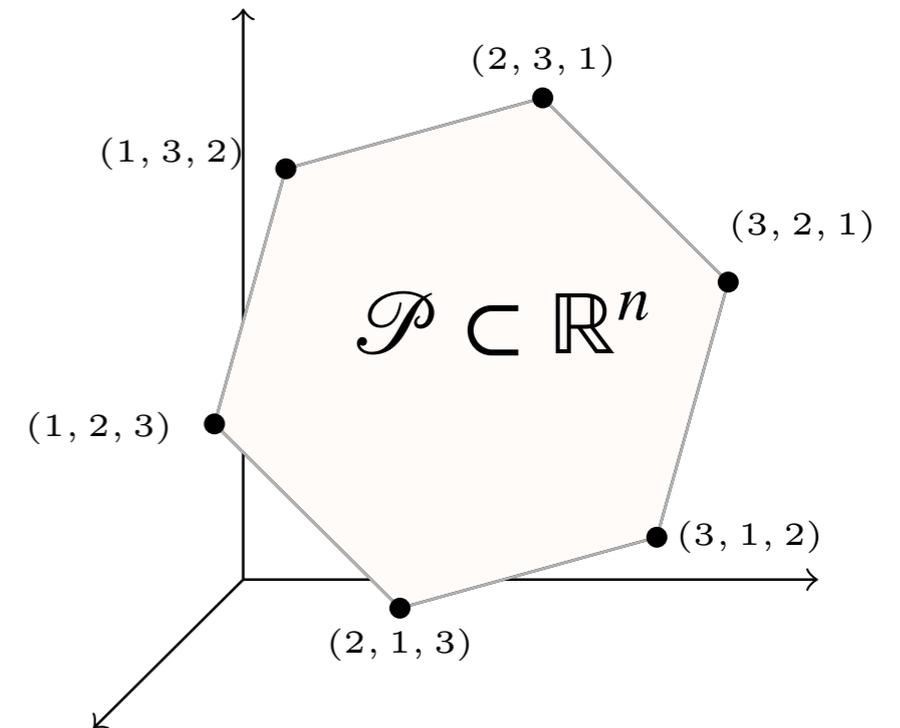


Entropy

Sinkhorn

This work

Permutahedron



L2 or Entropy

Pool Adjacent Violators (PAV)

Strategy outline

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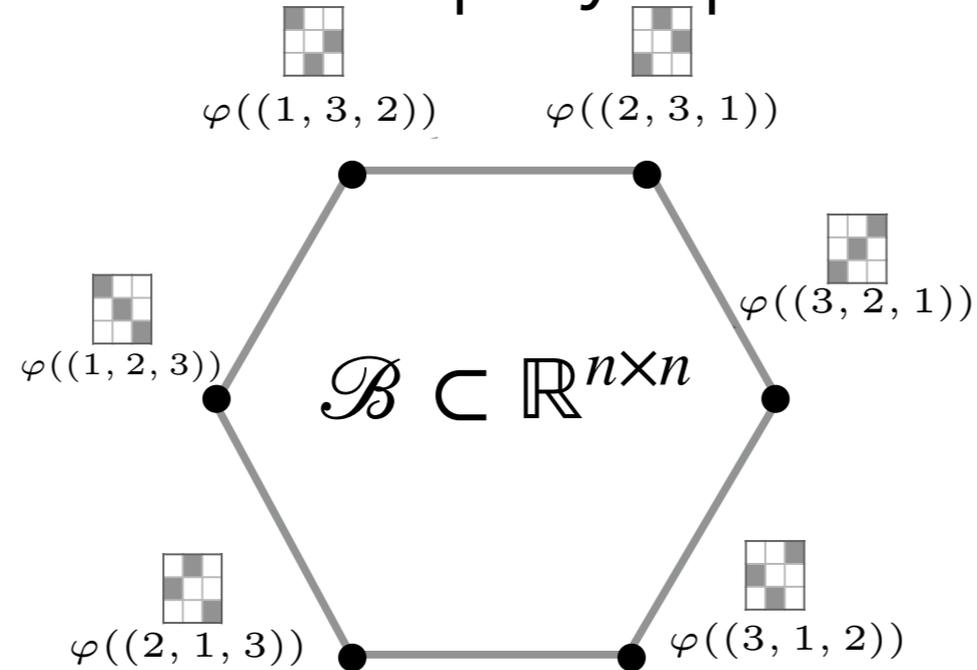
2. Regularization

3. Computation

4. Differentiation

Cuturi et al. [2019]

Birkhoff polytope



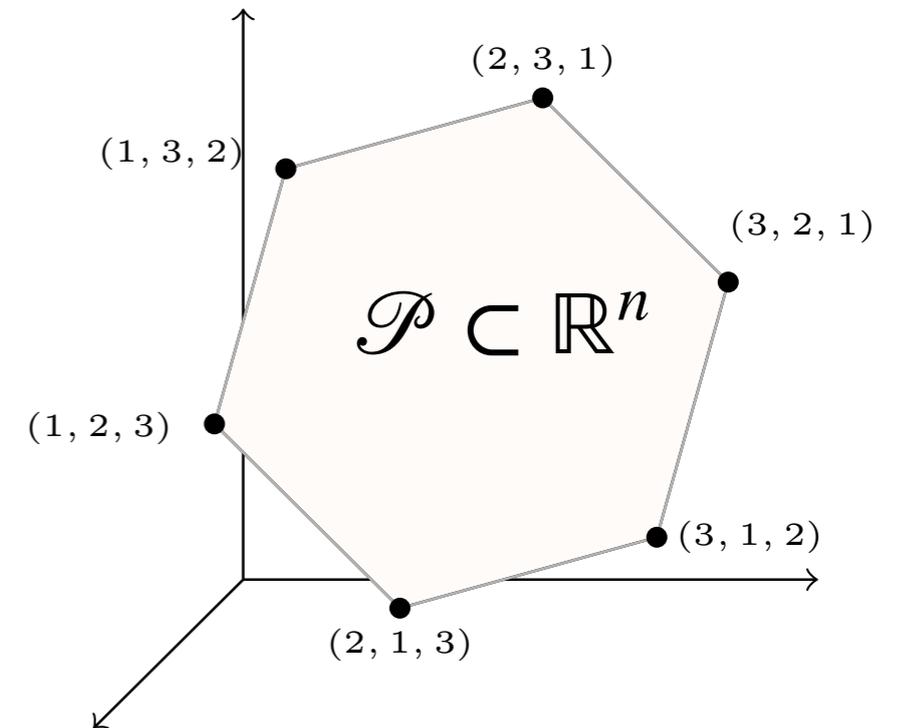
Entropy

Sinkhorn

Backprop through
Sinkhorn **iterates**

This work

Permutahedron



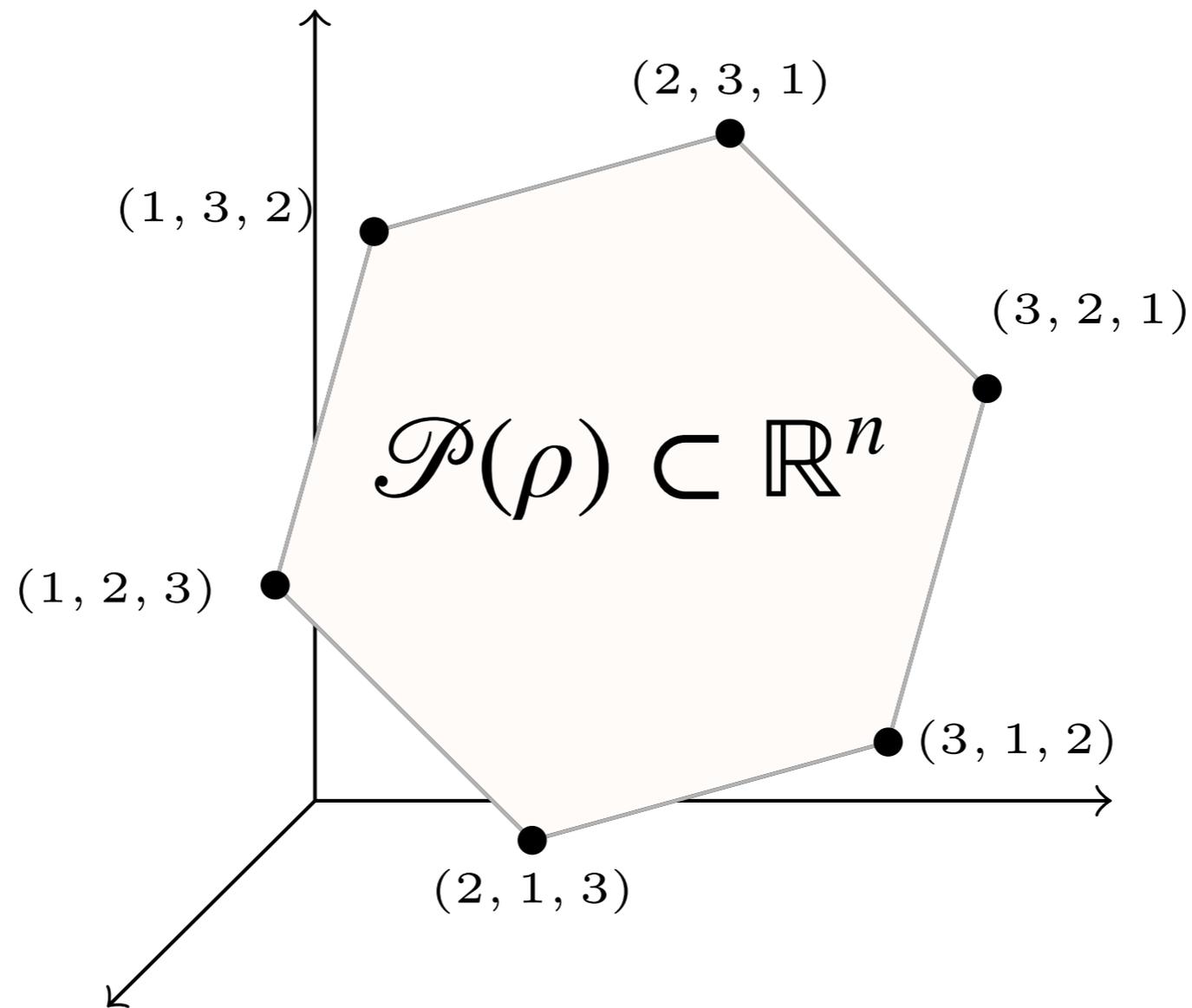
L2 or Entropy

Pool Adjacent
Violators (PAV)

Differentiate
PAV **solution**

Permutahedron

$$\mathcal{P}(w) \triangleq \text{conv}(\{w_\sigma : \sigma \in \Sigma\}) \subset \mathbb{R}^n$$



$$\rho \triangleq (n, n-1, \dots, 1)$$

Step 1: linear programming formulations

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Proposition

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$$s(\theta) = \arg \max_{y \in \mathcal{P}(\theta)} \langle y, \rho \rangle$$

$$r(\theta) = \arg \max_{y \in \mathcal{P}(\rho)} \langle y, -\theta \rangle$$

$$\rho \triangleq (n, n-1, \dots, 1)$$

Proof of the first claim

$$\rho_n > \rho_{n-1} > \dots > 1 \Rightarrow \sigma(\theta) = \arg \max_{\sigma \in \Sigma} \langle \theta_\sigma, \rho \rangle$$

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$$\begin{aligned} s(\theta) &\triangleq \theta_{\sigma(\theta)} \\ &= \arg \max_{\theta_\sigma: \sigma \in \Sigma} \langle \theta_\sigma, \rho \rangle \end{aligned}$$

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Step 2: introducing regularization

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Quadratic regularization $Q(y) \triangleq \frac{1}{2} \|y\|^2$

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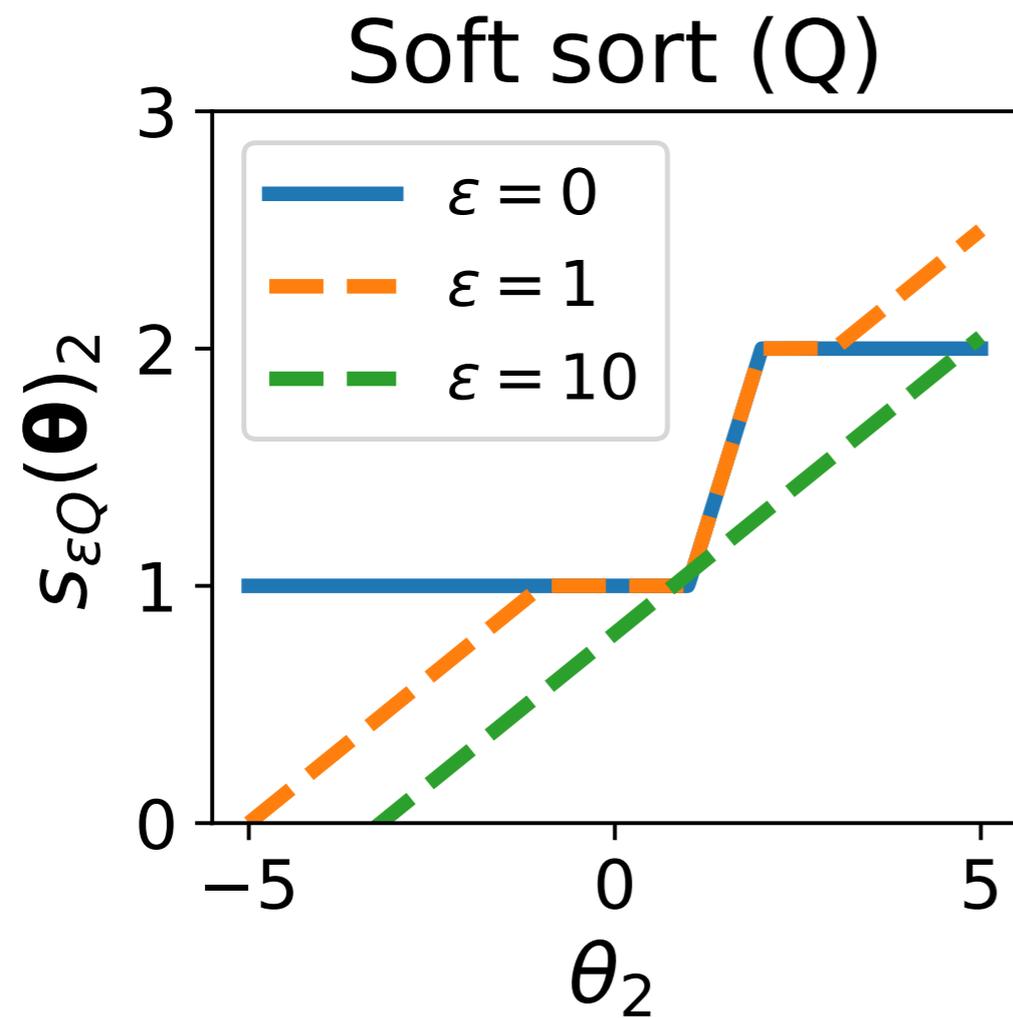
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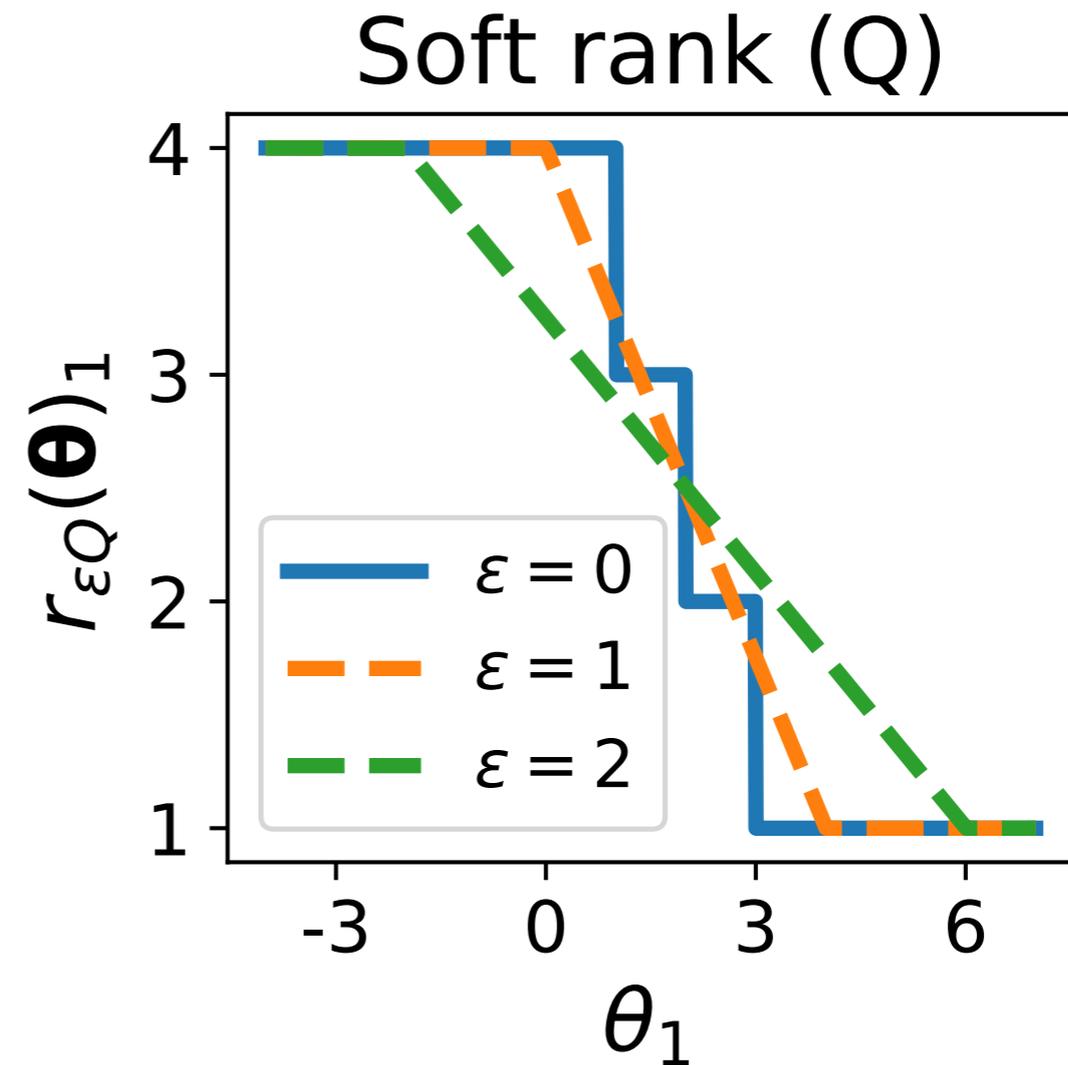
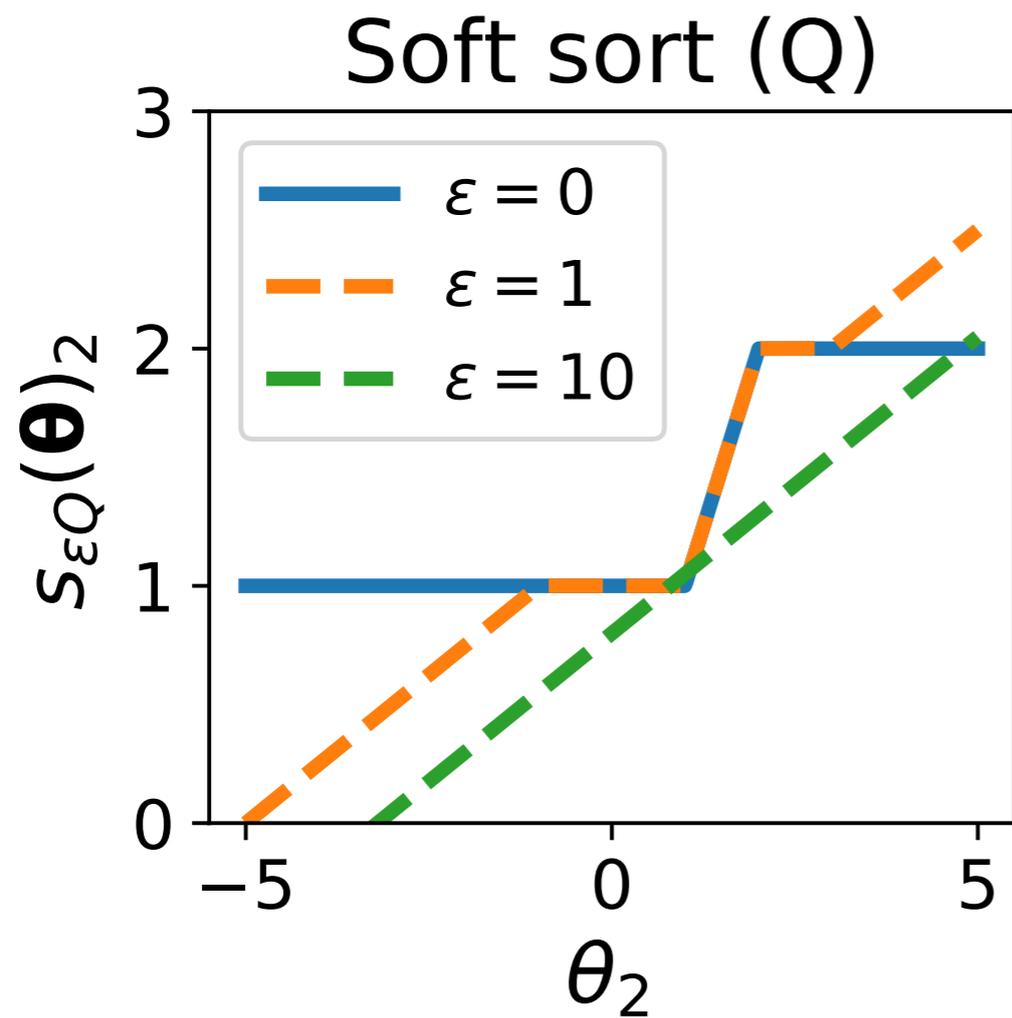
$$s_{\varepsilon Q}(\theta) \triangleq P_{\varepsilon Q}(\rho, \theta) = P_Q(\rho/\varepsilon, \theta)$$

$$r_{\varepsilon Q}(\theta) \triangleq P_{\varepsilon Q}(-\theta, \rho) = P_Q(-\theta/\varepsilon, \rho)$$

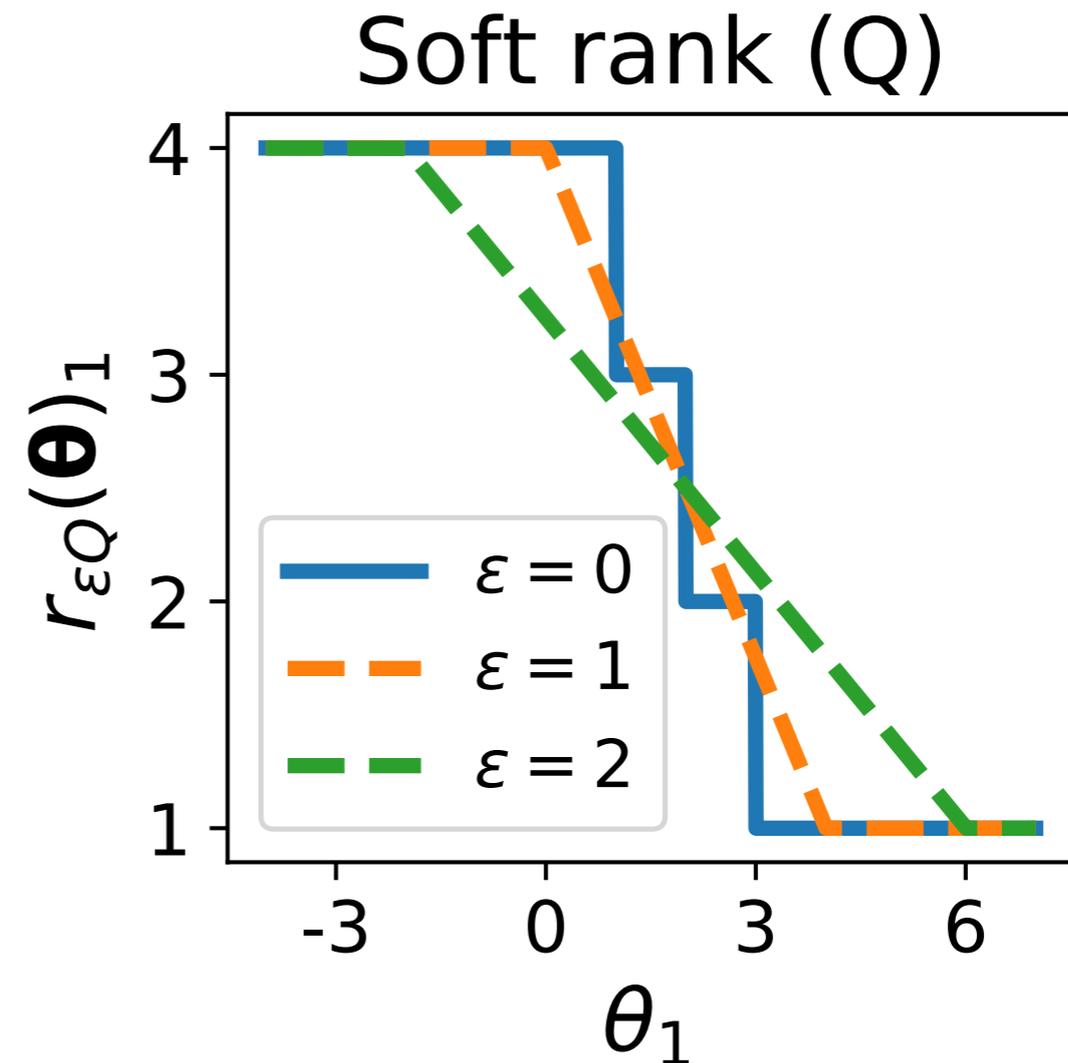
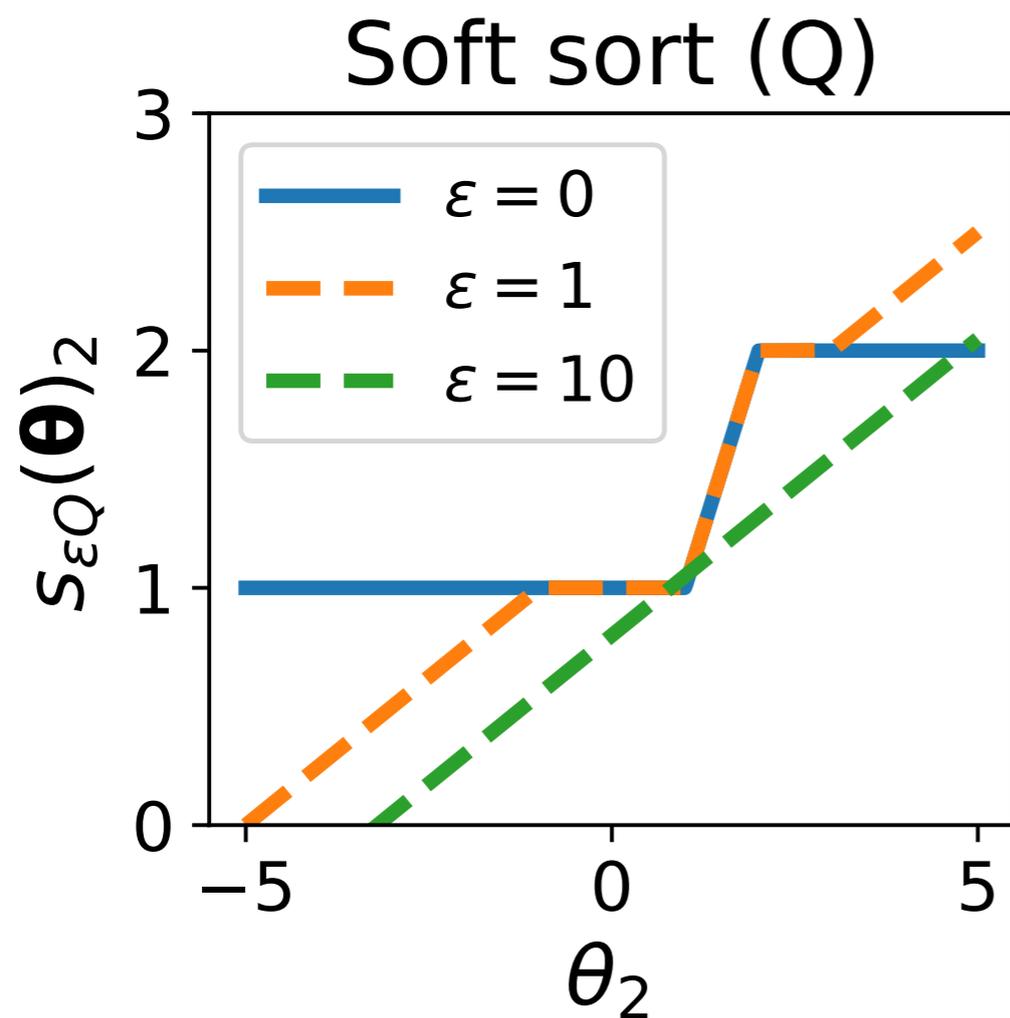
Continuity and differentiability



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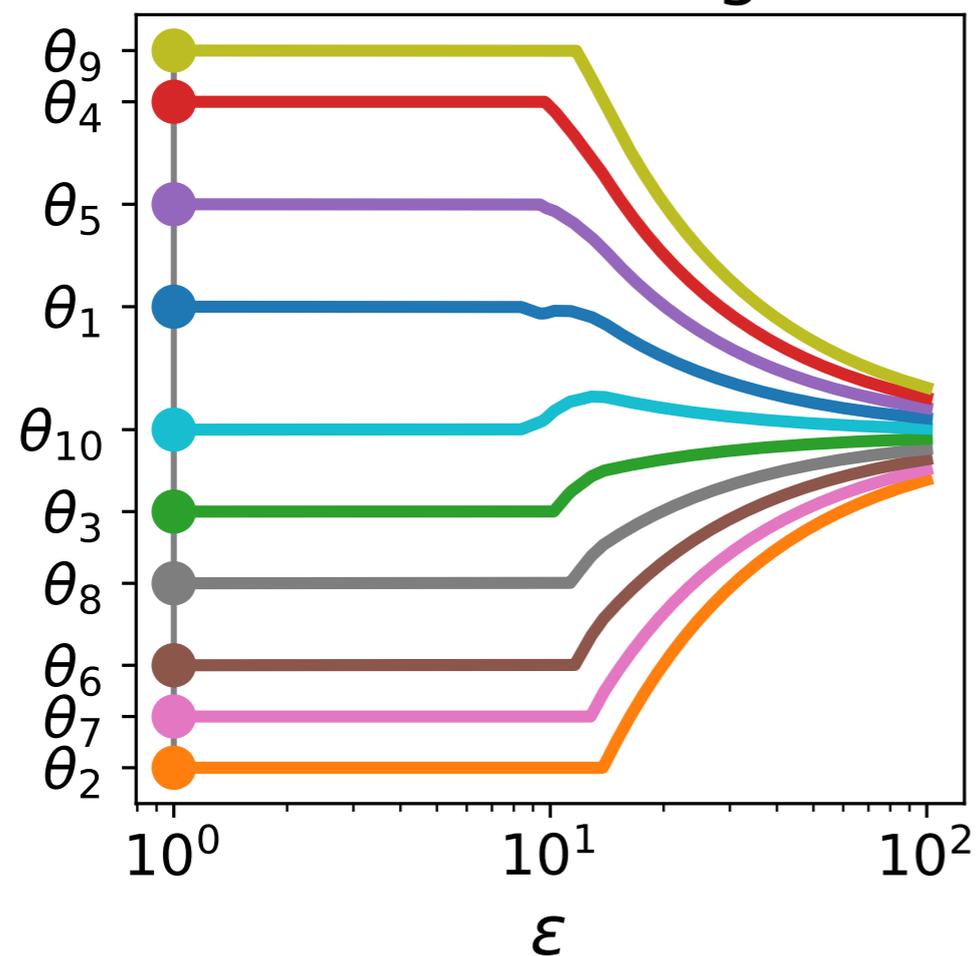


Properties

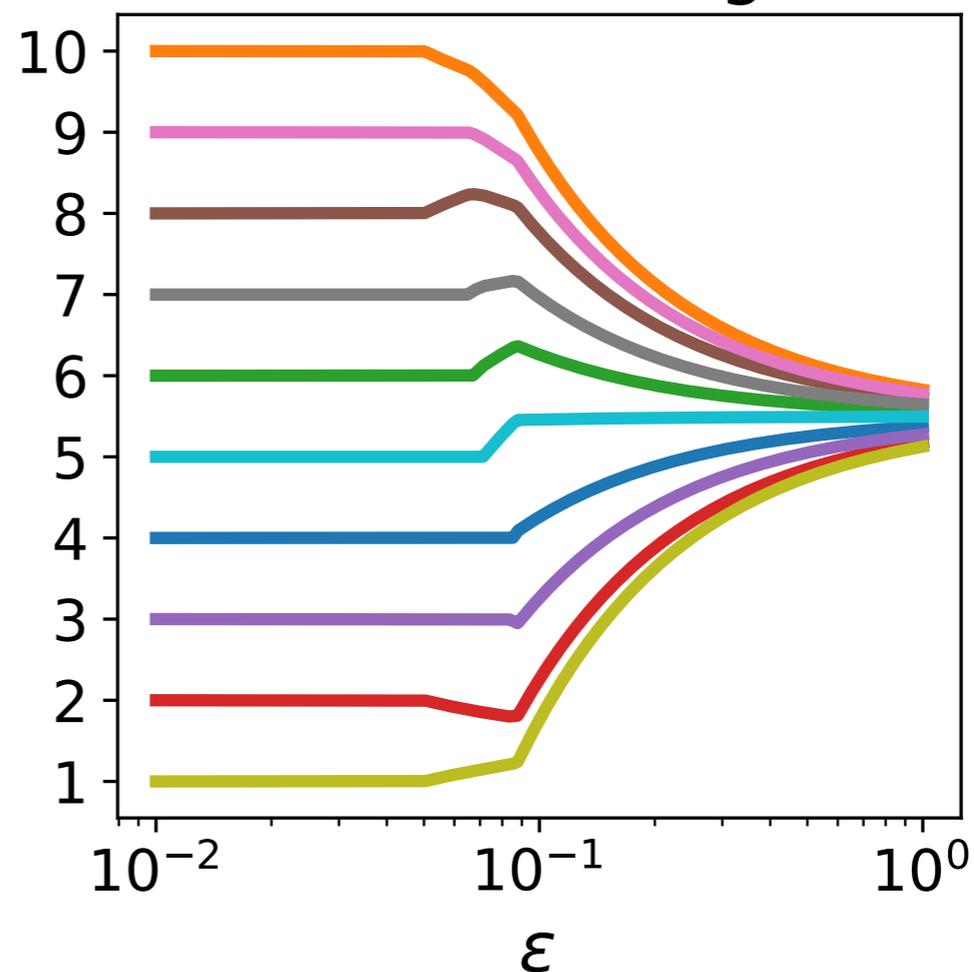
s_Q and r_Q are 1-Lipchitz continuous and differentiable almost everywhere.

Effect of regularization strength ε

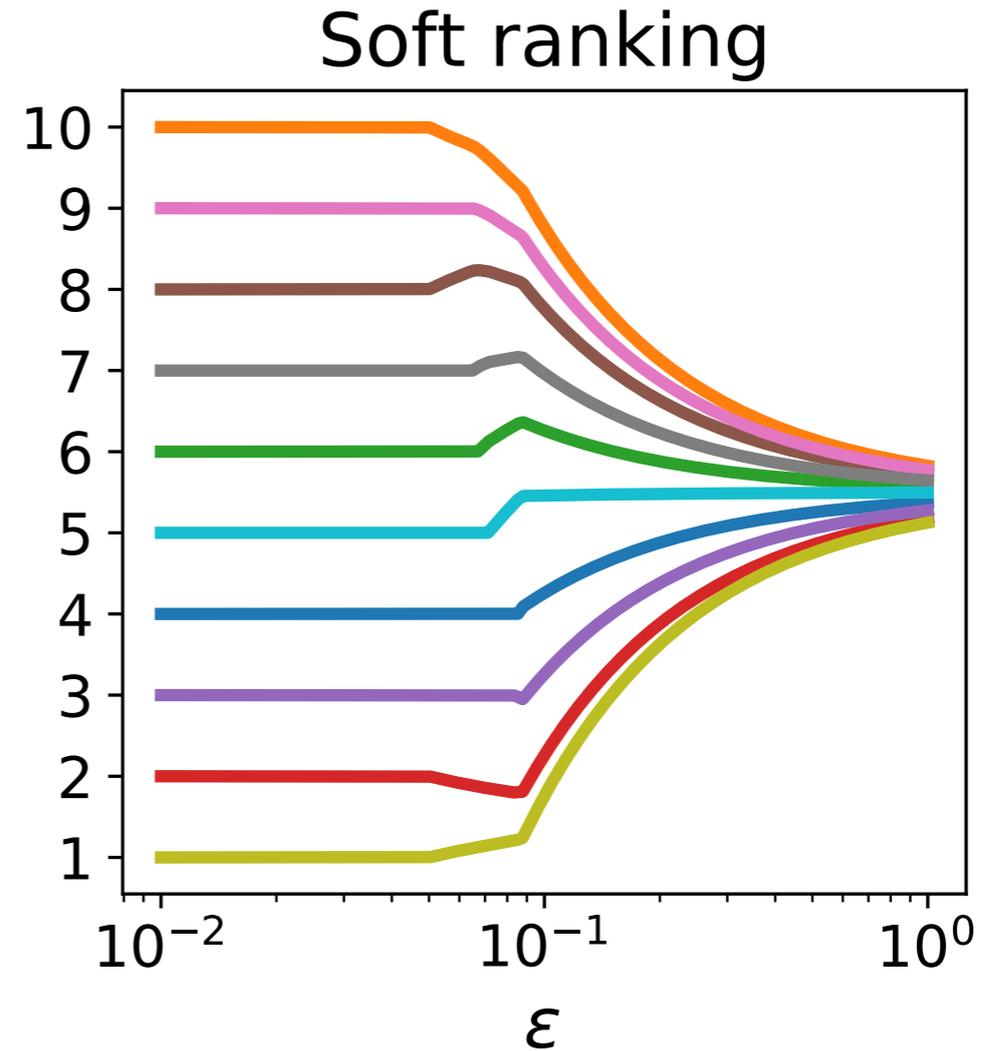
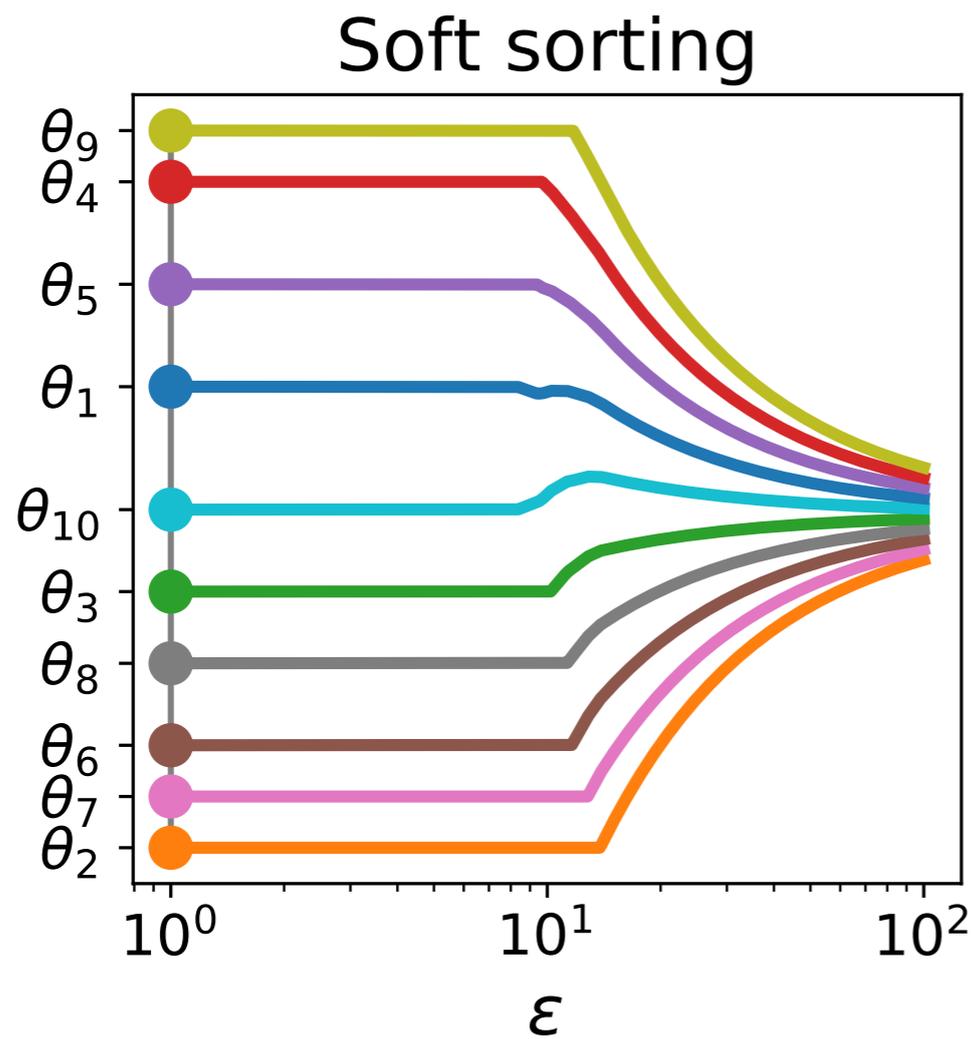
Soft sorting



Soft ranking



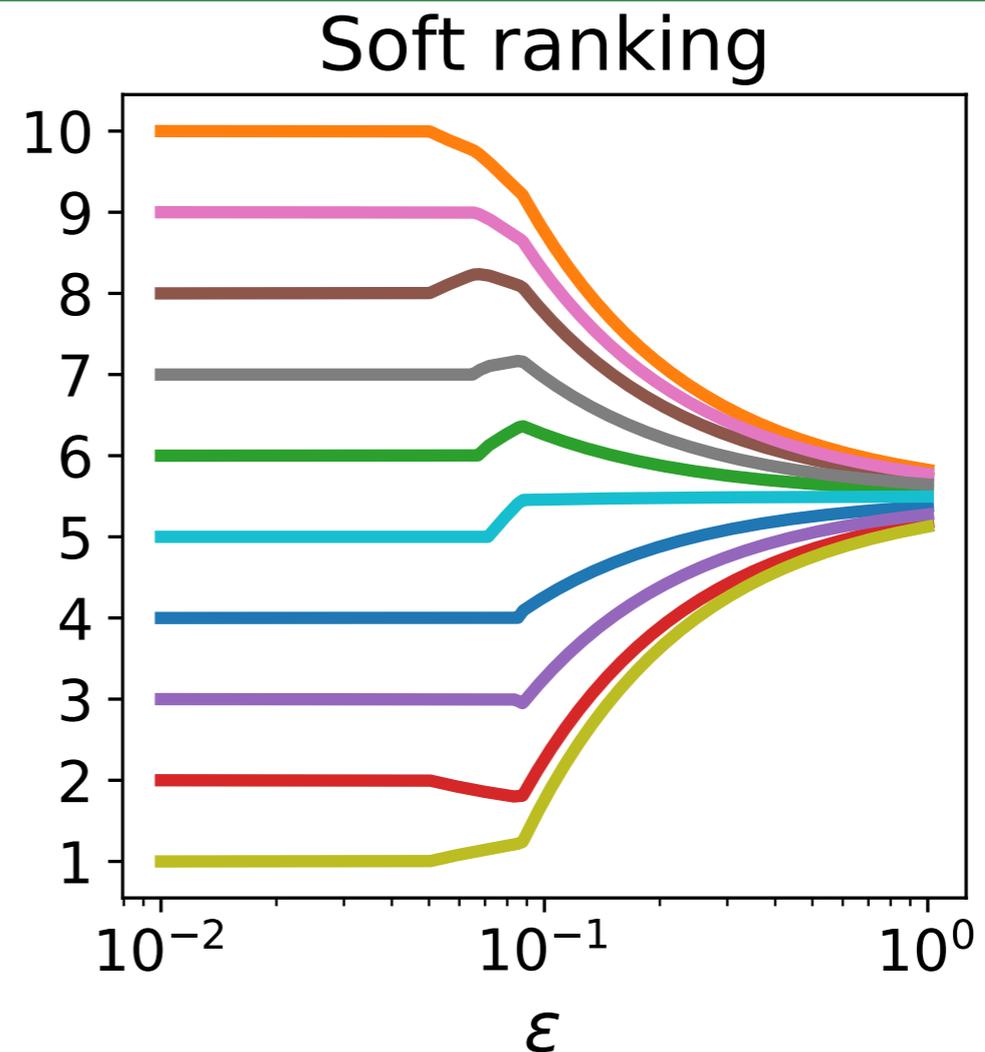
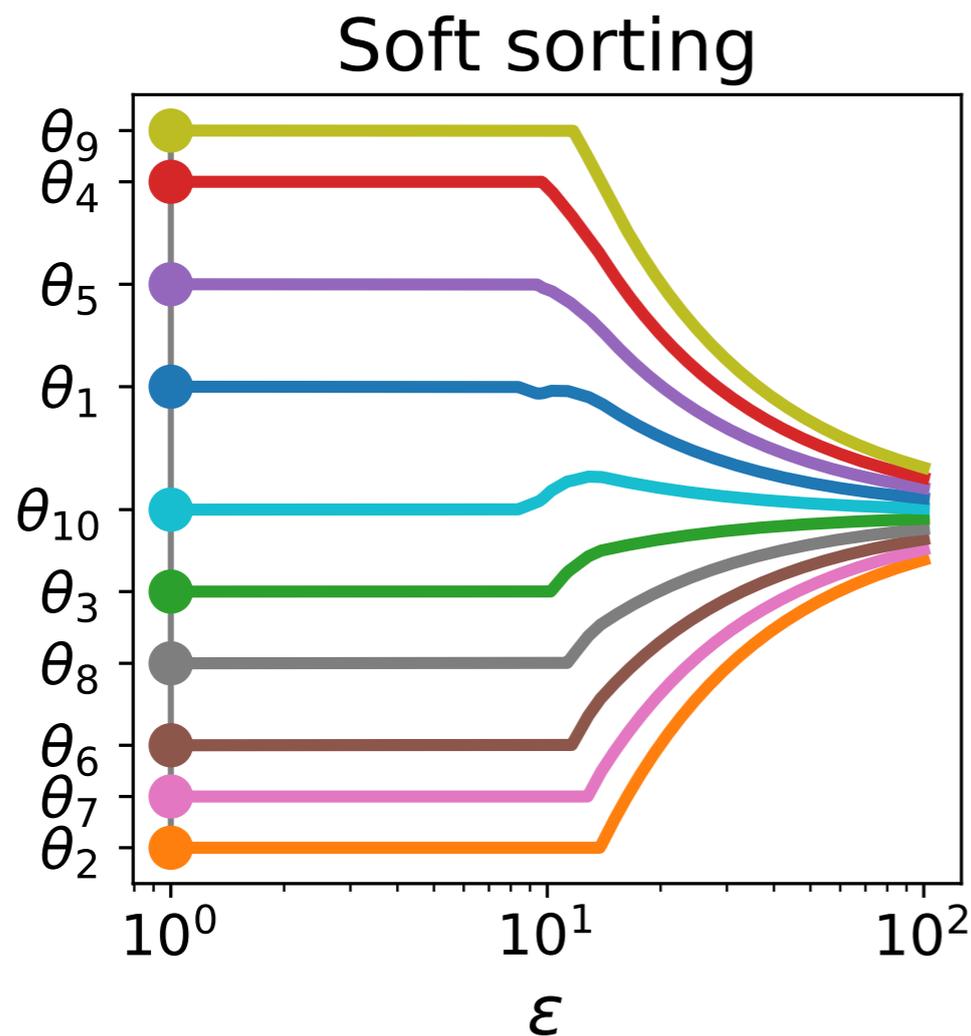
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Converge to hard version when $\varepsilon \leq \varepsilon_{min}$

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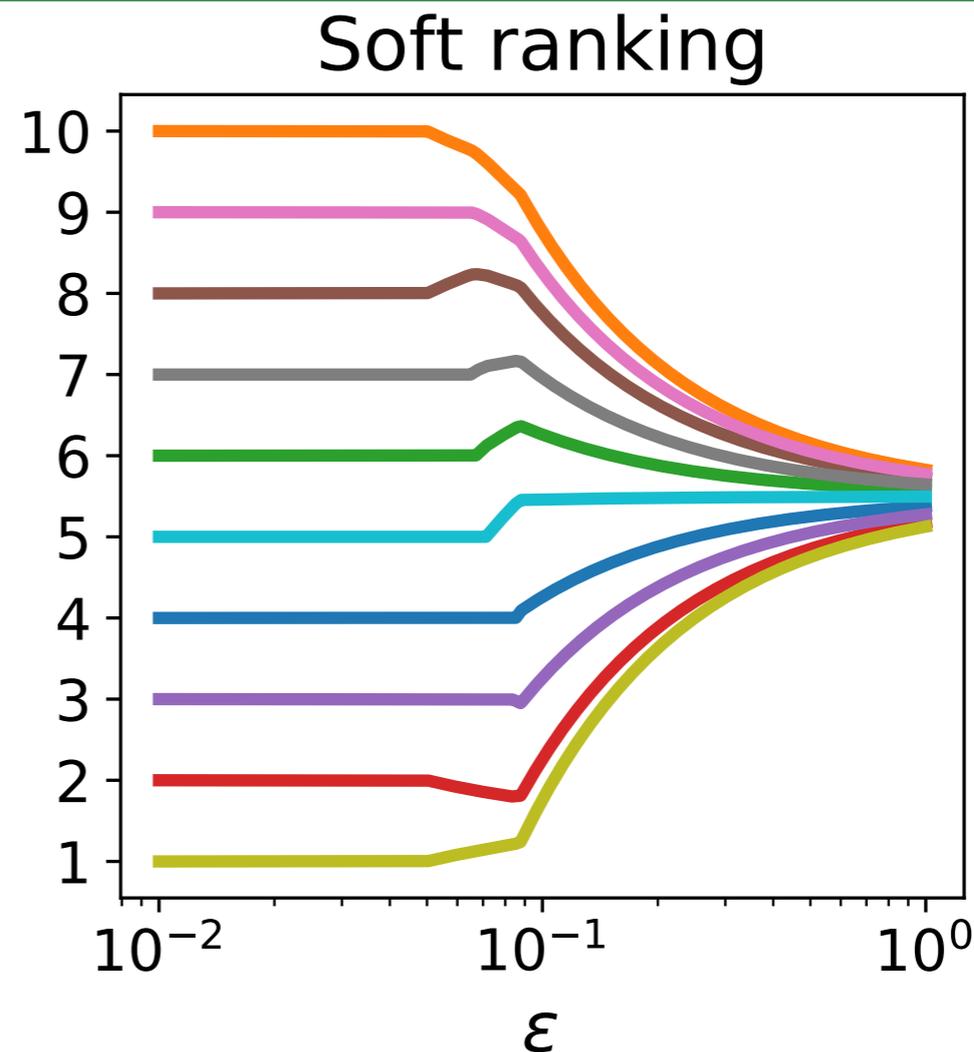
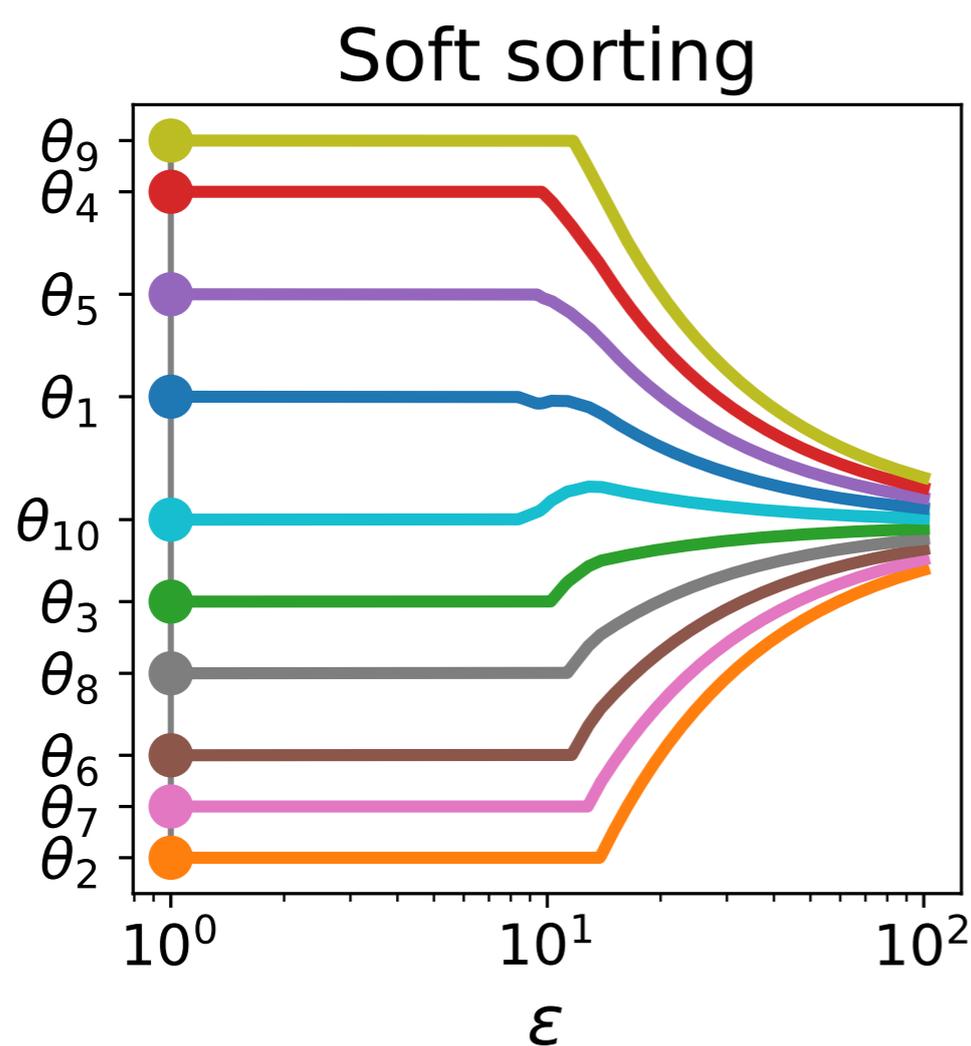


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Collapse to a mean when $\varepsilon \rightarrow \infty$

Effect of regularization strength ε



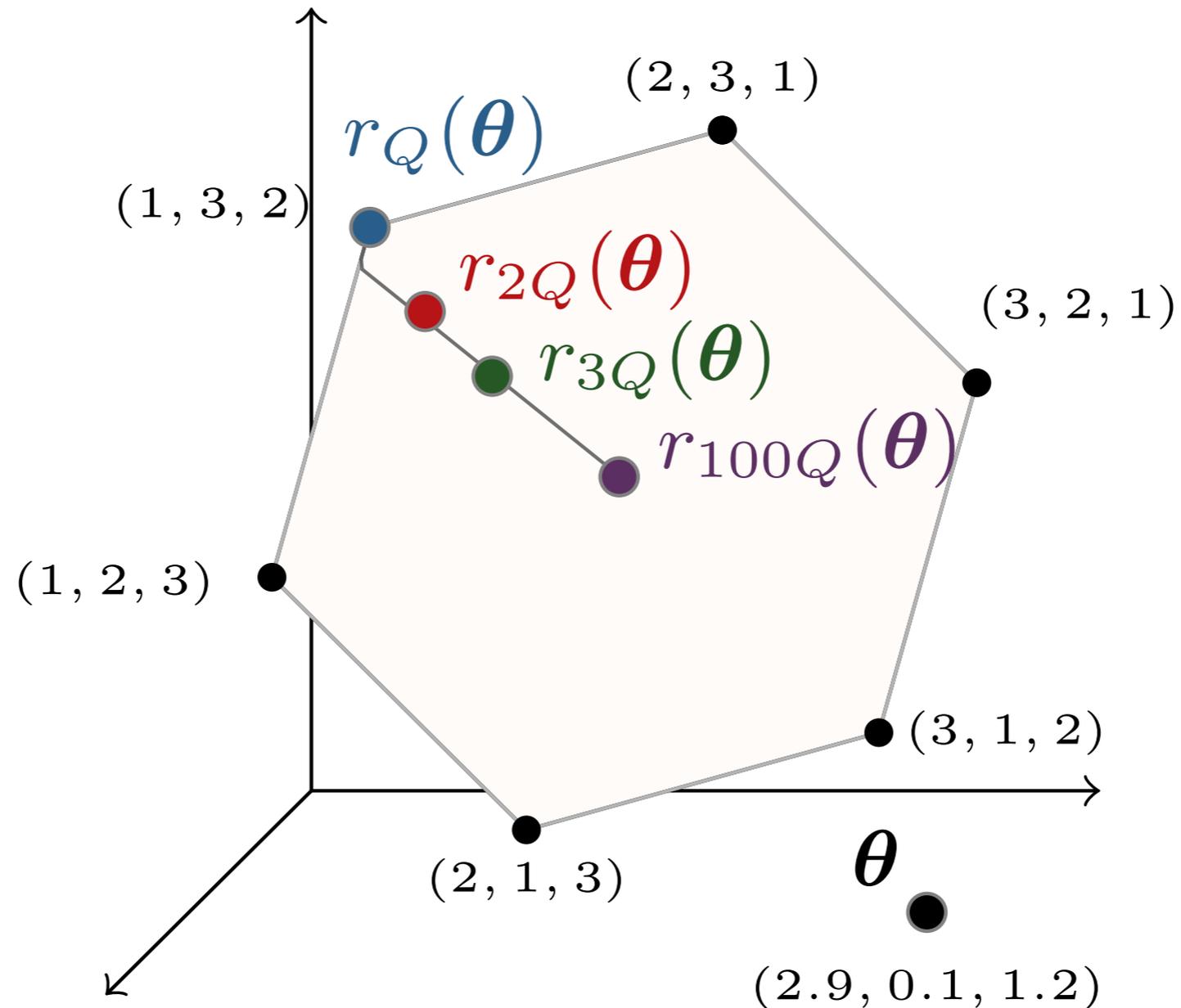
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Converge to hard version when $\varepsilon \leq \varepsilon_{min}$

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Order preserving (paths don't cross)

Regularization path



Collapse to a $\text{mean}(\rho)\mathbf{1}$ when $\varepsilon \rightarrow \infty$

Step 3: Computation

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Reduction to isotonic regression

Proposition

$$P_Q(z, w) = z - v_Q(z_{\sigma(z)}, w)_{\sigma^{-1}(z)}$$

$$v_Q(s, w) \triangleq \arg \min_{v_1 \geq \dots \geq v_n} \|v - (s - w)\|^2$$

Total time cost: $O(n \log n)$

e.g. [Negriño & Martins, 2014; Lim & Wright 2016]

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dual solution

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primal dual
relation

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Boils down to solving $v^{\star} = \arg \min_{v_1 \geq \dots \geq v_n} \|v - u\|^2$ $u = s - w$

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Pool Adjacent Violators (PAV): Finds a partition $\mathcal{B}_1, \dots, \mathcal{B}_m$ by repeatedly splitting coordinates. The worst-case cost is $O(n)$.

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Ex: $\mathcal{B}_1 = \{1,2\}$ $v_1^{\star} = v_2^{\star} = \text{mean}(u_1, u_2)$

$n=6$

$\mathcal{B}_2 = \{3\}$ $v_3^{\star} = \text{mean}(u_3) = u_3$

$\mathcal{B}_3 = \{4,5,6\}$ $v_4^{\star} = v_5^{\star} = v_6^{\star} = \text{mean}(u_4, u_5, u_6)$

Step 4: Differentiation

See also [Djologna & Krause, 2017]

Step 4: Differentiation

Differentiate $v_Q(s, w) = \arg \min_{v_1 \geq \dots \geq v_n} \|v - (s - w)\|^2$ w.r.t. s and w

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Differentiate $v_Q(s, w) = \arg \min_{v_1 \geq \dots \geq v_n} \|v - (s - w)\|^2$ w.r.t. s and w

Proposition

$$\frac{\partial v_Q(s, w)}{\partial s} = \begin{bmatrix} \mathbf{B}_1 & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \ddots & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{B}_m \end{bmatrix} \in \mathbb{R}^{n \times n}$$

$$\mathbf{B}_j \triangleq \mathbf{1} / |\mathcal{B}_j| \in \mathbf{R}^{|\mathcal{B}_j| \times |\mathcal{B}_j|}, \quad j \in [m]$$

Step 4: Differentiation

Differentiate $P_Q(z, w)$ w.r.t. z and w

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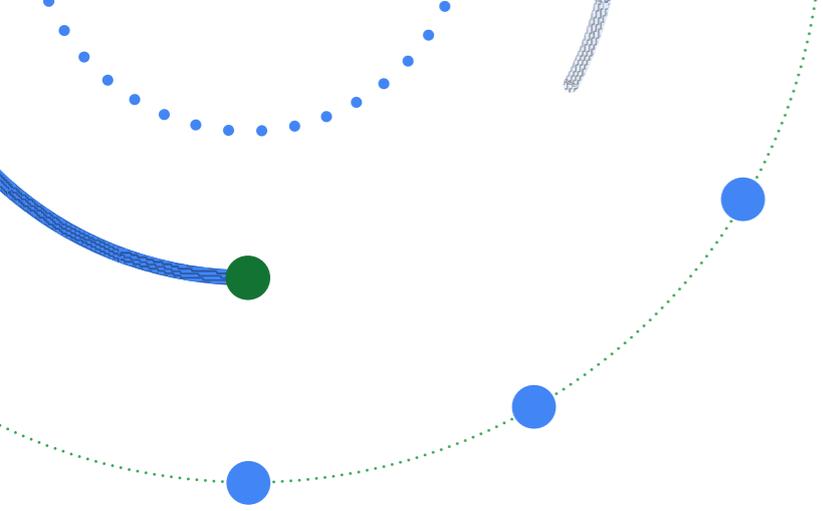
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$$\frac{\partial P_Q(z, w)}{\partial z} = J_Q(z_{\sigma(z)}, w)_{\sigma^{-1}(z)}$$

$$J_Q(s, w) \triangleq I - \frac{\partial v_Q(s, w)}{\partial s}$$

Multiplication with the Jacobian in $O(n)$ time and space (see paper)



Background

Proposed method

Experimental results

Robust regression

Robust regression

Least squares (LS)

$$\min_w \frac{1}{n} \sum_{i=1}^n \ell_i(w) \quad \ell_i(w) \stackrel{\triangle}{=} \frac{1}{2} (y_i - g_w(x_i))^2$$

ith loss

Robust regression

Least squares (LS)

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Soft Least trimmed squares (SLTS)

$$\min_w \frac{1}{n-k} \sum_{i=k+1}^n \ell_i^\varepsilon(w) \quad \ell_i^\varepsilon(w) \stackrel{\Delta}{=} [s_{\varepsilon Q}(\ell(w))]_i$$

ith "soft sorted" loss

Robust regression

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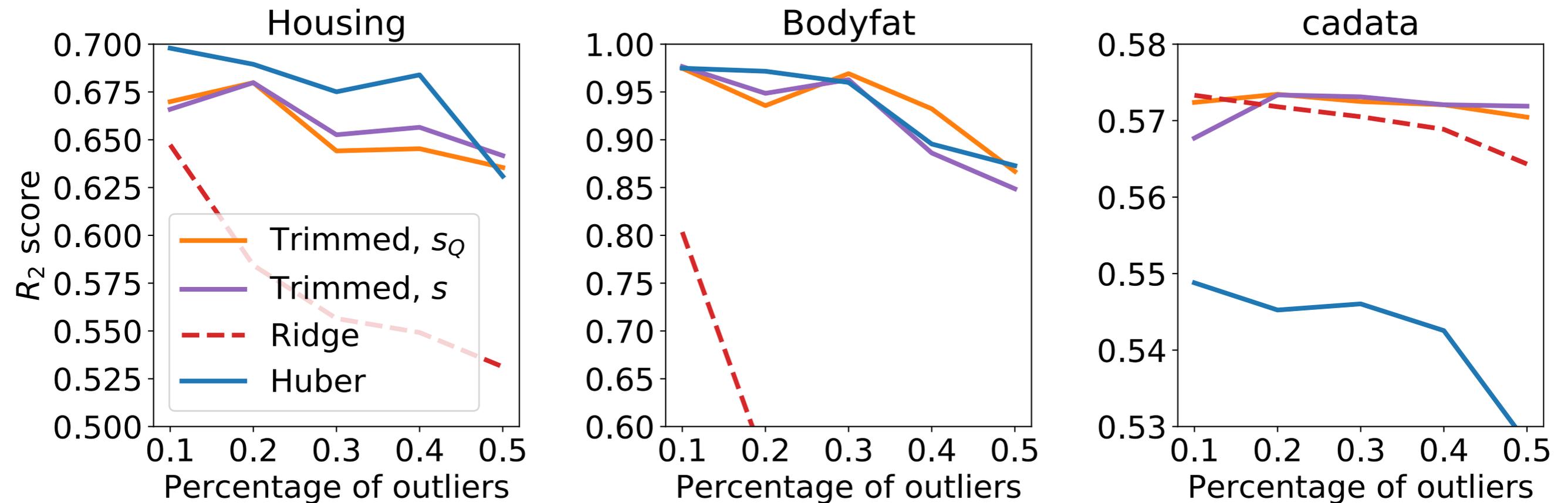
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$$\varepsilon \rightarrow 0 \quad SLTS \rightarrow LTS \quad \varepsilon \rightarrow \infty \quad SLTS \rightarrow LS$$

Robust regression



Evaluation: 10-fold CV

Hyper-parameter selection: 5-fold CV

Top-k classification

$$\ell : [n] \times \mathbb{R}^n \rightarrow \mathbb{R}_+ \quad \text{Cuturi et al. [2019]}$$

Ground truth
Predicted soft ranks

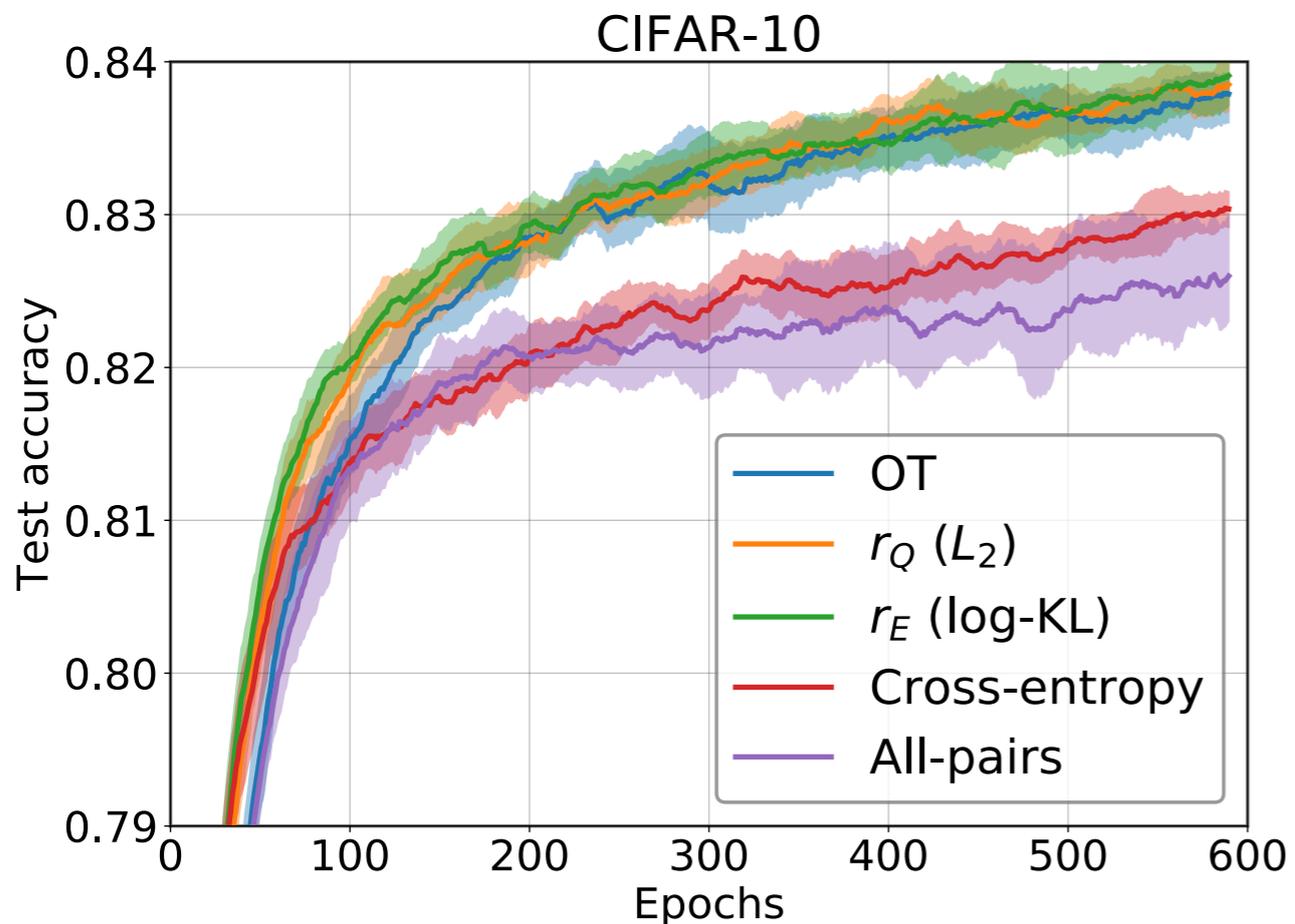
Top-k classification

$$\ell : [n] \times \mathbb{R}^n \rightarrow \mathbb{R}_+$$

Cuturi et al. [2019]

Ground truth
truth

Predicted
soft
ranks



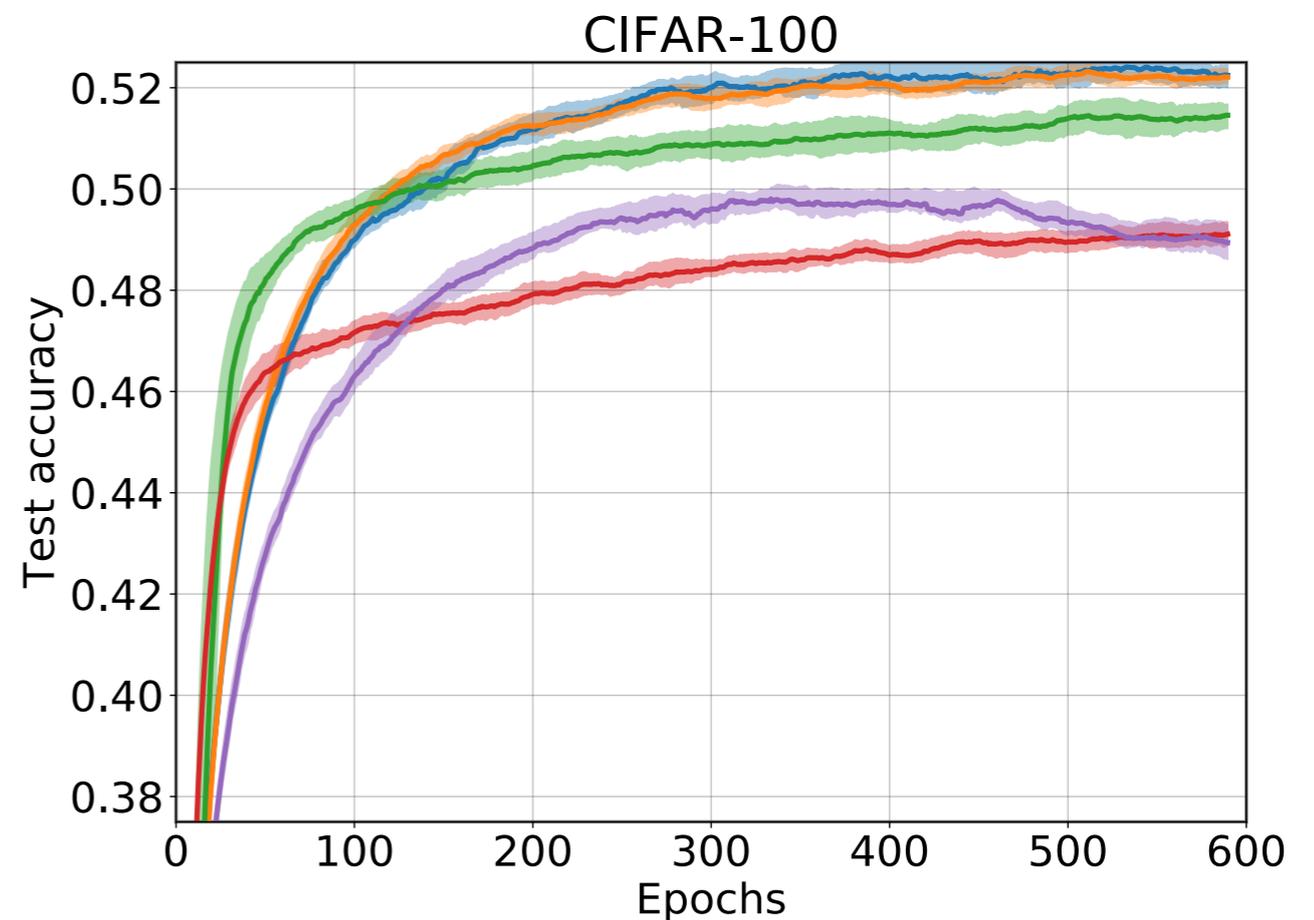
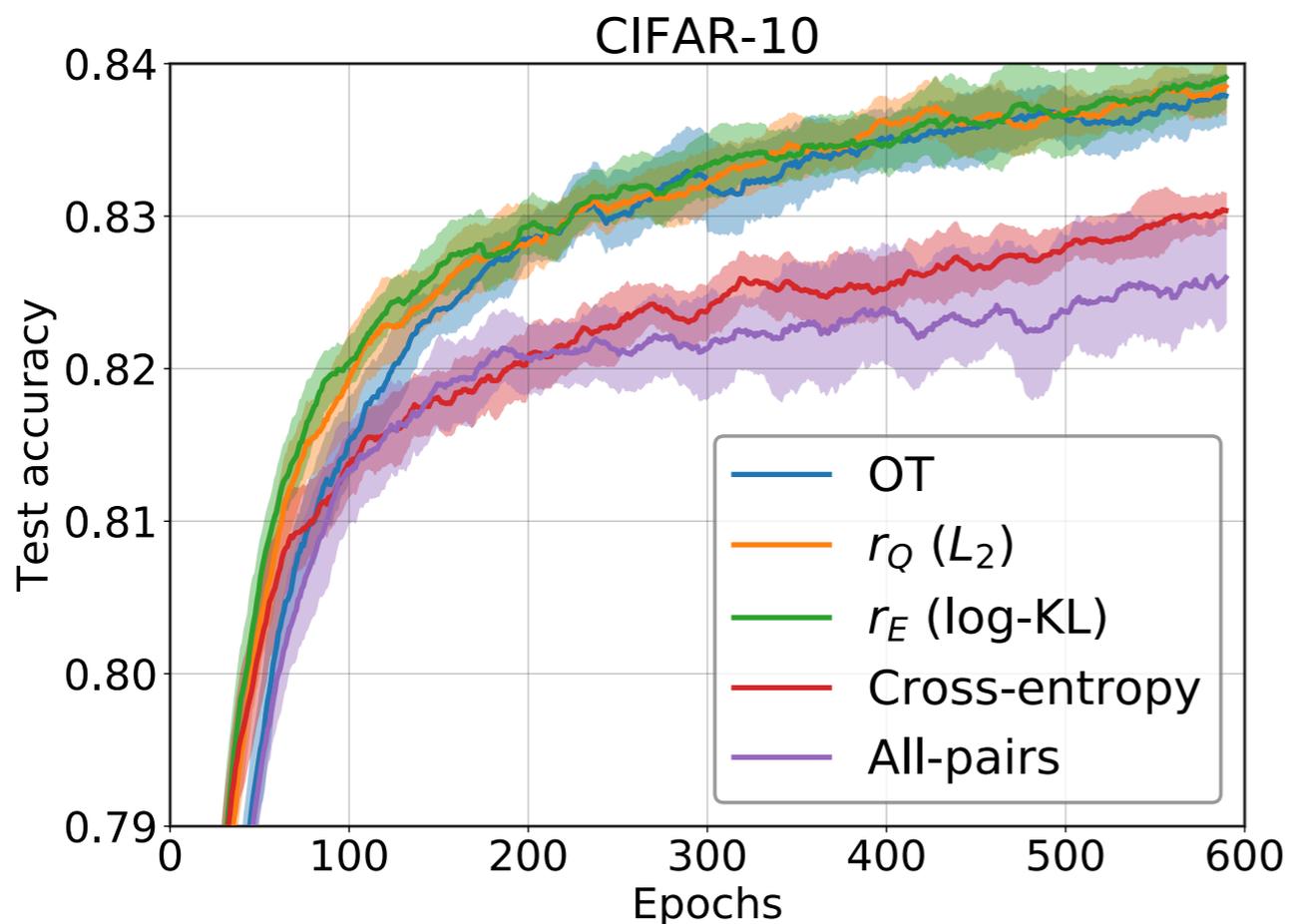
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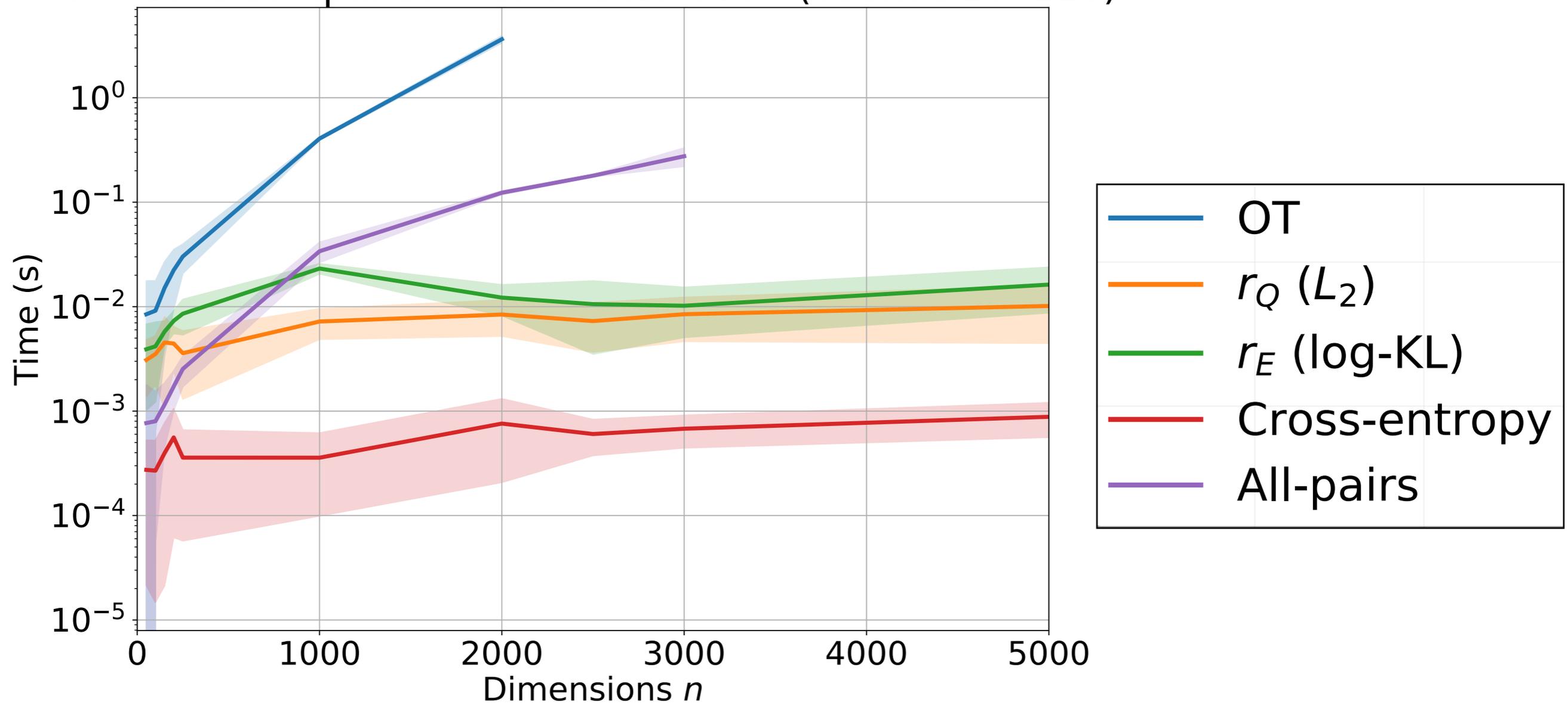
Ground truth
truth

Predicted
soft
ranks



Speed benchmark

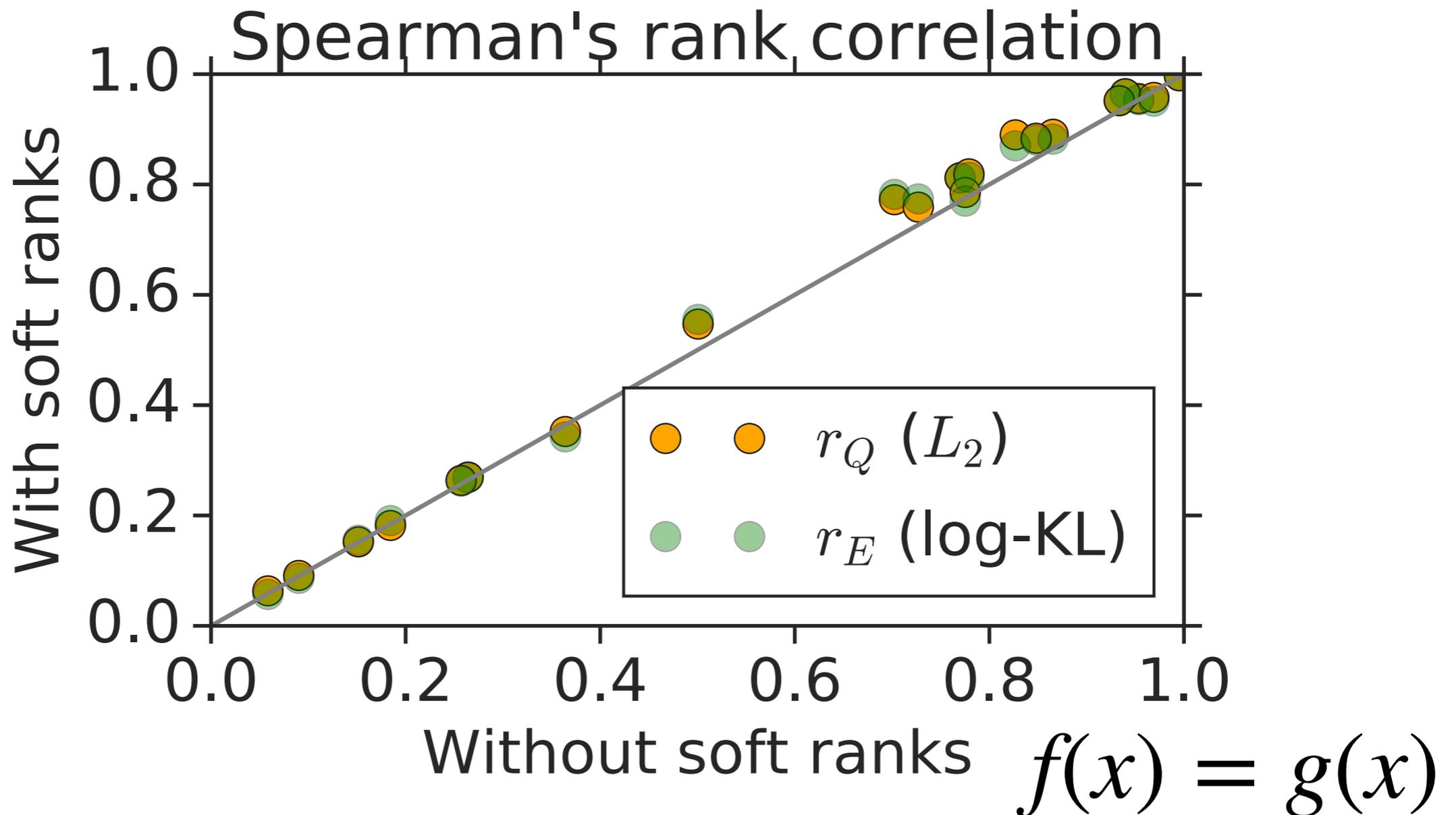
Runtime comparison for one iteration (batch size: 128)



Label ranking experiment

$$\ell_i \triangleq \frac{1}{2} \|y_i - f(x_i)\|^2 \quad y_i \in \Sigma$$

$$f(x) = r_Q(g(x))$$



Comparison on **21** datasets, 5-fold CV

Summary

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- Applications to least trimmed squares, top-k classification and label ranking

Preprint: Fast Differentiable Sorting and Ranking [arXiv:2002.08871]

Code: coming soon!